

Report (Universitetet; Bergen. Matematisk institutt)

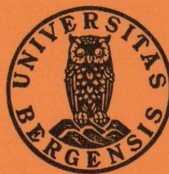
Department
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PURE MATHEMATICS

ON SHELLSORT AND THE FROBENIUS PROBLEM

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Abstract.

A bound $O(N^{1+1/2k})$ for the running time of Shellsort, with $O(\log N)$ passes, is proved very simply by application of a Frobenius basis with k elements.

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1. Shellsort theory.

For a description of Shellsort and the number-theoretical problem of Frobenius, we refer to two recent papers by Sedgewick [5] and by Incerpi and Sedgewick [3]. In the former paper, Sedgewick improves the Shellsort bound $O(N^{3/2})$ to $O(N^{1+1/2k})$, using a "result of Selmer" [7]. It is of course nice for a number-theoretician to see that his "useless" mathematics can really be applied. In all fairness, however, it should be made clear (as stated in [7]) that "my" result is really due to Hofmeister [3], as a special case of a general and rather complicated theorem. What I did in [7] was to give a direct, simple proof for this special case.

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Later (but published before [5]), Incerpi and Sedgewick [4] have improved the bound $O(N^{1+1/2k})$ to $O(N^{1+c})$, and further to $O(N^{1+c/\log k})$. In both cases, they circumvent the standard approach of Frobenius bases. Their proof of the latter bound is very nice, and I cannot in any way improve on it. Their proof of the bound $O(N^{1+c})$, does, however, result in an unnecessarily complicated increment sequence. The purpose of the present paper is to describe a simpler method, using a classical result in Frobenius theory.

In [4] p. 217, an increasing "base sequence" $\{a_i\} = a_1, a_2, \dots$ of natural numbers is used to produce the increments h_i of a Shellsort. A number c of different product sequences are interleaved, each such sequence consisting of certain products of c elements a_i . We shall see that one product sequence will suffice. In fact, we can define the increments by $h_i = 1$ and

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1. Shellsort theory.

For a description of Shellsort and the number-theoretical problem of Frobenius, we refer to two recent papers by Sedgewick [6] and by Incerpi and Sedgewick [4]. In the former paper, Sedgewick improves the Shellsort bound $O(N^{3/2})$ to $O(N^{4/3})$, using a "result of Selmer" [7] on a Frobenius basis with three elements. It is of course nice for a number-theoretician to see that his "useless" mathematics can really be applied. In all fairness, however, it should be made clear (as stated in [7]) that "my" result is really due to Hofmeister [3], as a special case of a general and rather complicated theorem. What I did in [7] was to give a direct, simple proof for this special case.

Later (but published before [6]), Incerpi and Sedgewick [4] have improved the bound $O(N^{4/3})$ to $O(N^{1+\epsilon})$, and further to $O(N^{1+\epsilon/\sqrt{\log N}})$. In both cases, they circumvent the standard approach of Frobenius bases. Their proof of the latter bound is very nice, and I cannot in any way improve on it. Their proof of the bound $O(N^{1+\epsilon})$ does, however, result in an unnecessarily complicated increment sequence. The purpose of the present paper is to describe a simpler method, using a classical result in Frobenius theory.

In [4] p. 217, an increasing "base sequence" $\{a_i\} = a_1, a_2, \dots$ of natural numbers is used to produce the increments h_j of a Shellsort. A number c of different product sequences are interleaved, each such sequence consisting of certain products of c elements a_i . We shall see that one product sequence will suffice. In fact, we can define the increments by $h_1 = 1$ and

$$h_j = a_{j-1} a_j \cdots a_{j+c-2}, \quad j > 1.$$

From these, we form a Frobenius basis with $c + 1$ elements:

$$(1) \quad B_{c+1}^{(j)} = \{h_{j+1}, h_{j+2}, \dots, h_{j+c+1}\} = \{b_1, b_2, \dots, b_{c+1}\}$$

(say). Under the condition (5) below for the base sequence $\{a_i\}$, we can then determine explicitly the Frobenius number

$$(2) \quad g(B_{c+1}^{(j)}) = \sum_{i=2}^{c+1} a_{j+i-2} h_{j+i} - \sum_{i=1}^{c+1} h_{j+i}.$$

This expression is clearly $O(h_j^{1+1/c})$ (if each term a_i is within a constant factor of the previous one). Just as in Theorem 2 of [4], we then get the running time of Shellsort bounded by $O(N^{1+1/(c+1)})$.

2. Frobenius theory.

We operate with a Frobenius basis

$$B_k = \{b_1, b_2, \dots, b_k\}, \quad \gcd(b_1, b_2, \dots, b_k) = 1.$$

Already Frobenius realized that a determination of $g(B_k)$ in the general case was extremely difficult. He therefore invited his audiences to look for good upper bounds for $g(B_k)$.

The first such bound was given already in the 1942 paper by A. Brauer [1] (indeed the first "serious" paper to be written on the problem of Frobenius). Let

$$d_0 = 0, \quad d_1 = b_1; \quad d_i = \gcd(b_1, b_2, \dots, b_i), \quad 2 \leq i \leq k.$$

Then Brauer showed that

$$(3) \quad g(B_k) \leq \sum_{i=1}^k b_i \left(\frac{d_{i-1}}{d_i} - 1 \right),$$

with equality if the following condition is satisfied:

$$(4) \quad \left\{ \begin{array}{l} \text{For all } i = 2, 3, \dots, k-1, \quad b_{i+1}/d_{i+1} \text{ is a linear} \\ \text{combination of } b_1/d_i, b_2/d_i, \dots, b_i/d_i \text{ with non-} \\ \text{negative integer coefficients.} \end{array} \right.$$

Further, Brauer and Seelbinder [2] showed that this condition is also necessary for equality in (3).

The proofs in the two papers quoted are rather complicated. Later, a very simple proof of the above results has been given by Rödseth [5].

We now apply this to the basis (1), and illustrate in the case $c = 3$, hence $k = 4$. If we write $\gcd(m,n) = (m,n)$, then

$$\begin{aligned} b_1 &= a_j a_{j+1} a_{j+2}, & b_2 &= a_{j+1} a_{j+2} a_{j+3} \\ b_3 &= a_{j+2} a_{j+3} a_{j+4}, & b_4 &= a_{j+3} a_{j+4} a_{j+5} \\ d_2 &= a_{j+1} a_{j+2} & \text{if } (a_j, a_{j+3}) &= 1 \\ d_3 &= a_{j+2} & \text{if also } (a_{j+1}, a_{j+3}) &= (a_{j+1}, a_{j+4}) = 1 \\ d_4 &= 1 & \text{if also } (a_{j+2}, a_{j+3}) &= (a_{j+2}, a_{j+4}) = (a_{j+2}, a_{j+5}) = 1. \end{aligned}$$

We will therefore assume that the base sequence $\{a_i\}$ satisfies $(a_i, a_{i+r}) = 1$ for $r = 1, 2, 3$ and all $i \geq 1$. In the general case, the corresponding condition is

$$(5) \quad \gcd(a_i, a_{i+r}) = 1, \quad r = 1, 2, \dots, c, \quad i = 1, 2, \dots.$$

The conditions (4) are trivially satisfied, since

$$\frac{b_3}{d_3} = a_{j+3} a_{j+4} = a_{j+4} \frac{b_2}{d_2}, \quad \frac{b_4}{d_4} = a_{j+3} a_{j+4} a_{j+5} = a_{j+5} \frac{b_3}{d_3}.$$

In the general case, we similarly have

$$\frac{b_{i+1}}{d_{i+1}} = a_{j+c+i-1} \frac{b_i}{d_i}, \quad i = 2, 3, \dots, c.$$

We can thus use (3) with equality. Since $d_0 = 0$ and $d_{i-1}/d_i = a_{j+i-2}$, $i = 2, 3, 4$, we get

$$g(B_4^{(j)}) = \sum_{i=2}^4 a_{j+i-2} b_i - \sum_{i=1}^4 b_i,$$

where $b_i = h_{j+i}$. The generalization to (2) is immediate.

It remains to find base sequences $\{a_i\}$ satisfying (5). One obvious possibility is suggested in [4]: Choose $\alpha > 1$, and a_i as the smallest prime $\geq \alpha^i$.

An interesting alternative stems from Sedgewick's first paper [6], where his Theorem 6 in fact corresponds to $c = 2$ above (but he does not give $g(B_3^{(j)})$ explicitly), with

$$(6) \quad a_i = 2^{i+1} - 3 .$$

The conditions (5) for $r = 1, 2$ are clearly satisfied, since $a_{i+1} - a_i = 2^{i+1}$, $a_{i+2} - a_i = 3 \cdot 2^{i+1}$. In fact, the choice (6) is possible also for $c = 3$, since now $a_{i+3} - a_i = 7 \cdot 2^{i+1}$, and $7 \nmid a_i$. This stems from the fact that 2 is not a primitive root of 7, $2^3 \equiv 1$, and $2^t \not\equiv 3 \pmod{7}$ for all t .

For $c = 4$, we similarly form $a_{i+4} - a_i = (2^4 - 1)2^{i+1} = 3 \cdot 5 \cdot 2^{i+1}$. Since 2 is a primitive root both of 3 and of 5, we must now try to make $a_i = 2^{i+m} - n$, where n (odd) is divisible by 15. The smallest choice of n also possible modulo 7 is $n = 45$, so we can put

$$(7) \quad a_i = 2^{i+5} - 45 .$$

Quite surprisingly, this choice is possible for all $c \leq 9$, since

$$2^5 - 1 = 31, 2^6 - 1 = 3^2 \cdot 7, 2^7 - 1 = 127, 2^8 - 1 = 3 \cdot 5 \cdot 17, 2^9 - 1 = 7 \cdot 73 .$$

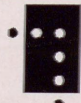
Here 2 is not a primitive root modulo any of the primes $p = 17, 31, 73, 127$, and it turns out that always $2^t \not\equiv 45 \pmod{p}$ for these four primes.

For $c = 10$, however, we have $2^{10} - 1 = 3 \cdot 11 \cdot 31$, where 2 is a primitive root of 11, so we must have $11 \mid n$. Trying to combine with the earlier primes considered, we end up with quite a large n .

For all sorting purposes, it is hardly practical to choose $c > 9$. We can then use (7), or (6) for $c = 2, 3$. I leave it to the sorting specialists (like Sedgewick) to test whether my above procedure for Shellsort can compete with procedures described earlier.

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