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Some Computing Aspects of Projective Geometry I.

Basic functions, algorithms and procedures.

By

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It therefore is appearent that the present project can not be carried out exclusively within the environment of symbolic manipulation. Nevertheless, the size of MACSYMA has been indispensable, and tools of this kine will have to be used extensively in the future. In fact, aside from the obvious epproach of trying to implement functions

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References:

Introduction.

The aim of this paper is to lay the foundation for certain investigations in projective geometry by means of a computer.

Here I only give some of the basic techniques which will have to play a central role in these investigations, to appear later as [Hm 8, 9]. Thus the present work is of a preliminary nature.

One of the main points at this stage has been to investigate the feasability of some of those basic techniques, implemented on a computer. For this I have been able to use the MACSYMA-system of the Mathlab-group at the Computer Science Laboratory of MIT, Cambridge, Mass. MACSYMA is a very sophisticated system for symbolic manipulation, in principle very well suited for the kind of investigation I am undertaking here and in the articles announced above.

However, there is a difficulty which became serious already at the current preliminary stage: Namely, as the size of the problem grows - for instance the number of transcendentals in the expressions one works with - then the size of the computations in some cases tend to grow exponentially. And since the elegant and flexible facilities available in MACSYMA also tend to require considerable core memory space, one frequently finds that the swell in intermediate calculations severely limits the size of the problems which can be treated. This is a serious obstacle. Moreover, for similar reasons some of the computations would tend to run for an unexpectedly long time, even in relatively simple cases.

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It therefore is appearent that the present project can not be carried out exclusively within the environment of symbolic manipulation. Nevertheless, the use of MACSIMA has been indispensable, and tools of this kind will have to be used extensively in the future. In fact, aside from the obvious approach of trying to implement functions written in a language for symbolic manipulation in a lower level language, one may develop programs by obtaining intermediate results for instance by MACSYMA and use these in programs written in a lower level language. A typical example of this is given by the expressions RE[1], RE[2] and RE[3] which have been found using MACSYMA. RE[3] could simply not have been computed "by hand", while RE[1] is easy. These expressions make it possible to find the Chern- and Segre classes for Grassmanians of lines, of planes and of 3-spaces in $\mathbb{P}^{\mathbb{N}}$ for all N, using conventional programming languages. Unfortunately MACSYMA was not able to compute RE[4]. See sections 2 and 5 for details on this.

Moreover, in my opinion there is no question that computing with MACSYMA competes favorably with traditional approaches such as the one taken by <u>A. Lascoux</u> in [Lx] to compute Chern classes of tensor products and the 2 nd. symmetric and exterior power: Indeed, in section 1 it is shown how this goes through smoothly for tensor products, which is used later. Also the second exterior power is treated as an example, and it is clear how to generalize this to any exterior power. The symmetric powers are dealt with similarly, but this is omitted here. From a computational point of view I believe that the present approach is preferable.

Also, it is instructive to compare <u>R. Donagi's</u> computations in [Do], the appendix, to the material in section 3.

Thus in section 1, I give a procedure^{*)} for computing the Chern polynomial of a tensor product in terms of the Chern polynomials of the factors. Using the procedure, a function carrying out the computation is then written in MACSYMA. In section 2, procedures for the Chern and Segre classes of Grassmanians are similarly given and implemented on MACSYMA. In section 3 an algorithm is given which will

*) The word "procedure" is used here in a less technical sense than the word "algorithm".

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convert expressions in the Chern classes of the universal quotient bundle of a Grassmanian into expressions in the Schubert symbols. This is done by repeatedly applying Pieri's formula. The algorithm is implemented on MACSYMA, but unfortunately the computations tend to be rather time-consuming.

In section 4 I give some basic formulae and functions for computing the embedding- and duality properties of projective varieties. This material should be viewed in light of section 6, where the continuation of this work is outlined.

Section 5 contains some combinatorical aspects of Grassmanians of lines, which also illustrates how parts of this project can be carried out with conventional computer programs.

I would like to thank Professor Joel Moses of the Computer Science Laboratory at MIT for giving me access to the MACSYMA system, and the entire Mathlab group for their patience and help during my work with MACSYMA. In particular I would like to thank Dr. B. Trager, whith whoom I had many enlightening conversations.

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- V1 -

§ 1. Computing the Chern classes of a tensor product.

One of the techniques which we will need in this paper, is a practicable algorithm for computation of the Chern classes of a tensor product $E \otimes F$ of two locally free 0_X - Modules on a scheme X, in terms of the Chern classes of E and F.

The need for such a method, as well as the related one for symmetric and exterior powers, arises in many situations. In addition to the questions studied in this paper, they are also needed for the higher order Thorn-Boardman singularities.

This is the motivation for a recent article by <u>A. Laxoux</u>, [Lx] where the theory of Schur-functions is utilized to obtain explicit formulae for these Chern classed.

However, while Laxoux's expressions are nice to have, and of some interest in their own right, they do not lend themselves easily to computation: The Chern classed in question are obtained in terms of certain determinants in the Segre classes of E and F, and even though the passage from Chern polynomials to Segre polynomials is trivial on a computer, the further computations with Laxoux's formulas would still be rather large.

Here we take a different and much simpler approach to this problem: Using only substitutions, expansions and simplifications, as well as the function RESULTANT, we write a function in MACSYMA,

TENSOR $(\arg_1, \arg_2, \arg_3, \arg_4)$, which when given the arguments

 $\arg_1 = 1 + c_1(E)T + \dots + c_e(E)T^e$,

- 1 -

6 1. Computing the Chern classes of a tensor product.

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TENSOR (are, areg, areg, areg, areg),

T(S), 0 + ... + T(S), 0 + 1 - 1378

- 1

the Chern polynomial of E;

 $\arg_2 = e$,

the rank of E;

$$\arg_{z} = 1 + c_{1}(F)T + \dots c_{f}(E)T^{I}$$
,

the Chern polynomial of F;

$$\arg_{\mu} = f$$
,

the rank of f; will return the Chern polynomial of $E \bigotimes F$ in terms of the Chern classes

$$c_1(E)$$
, ..., $c_e(E)$; $c_1(F)$, ..., $c_f(E)$.

The indeterminate must be T in \arg_1 and \arg_2 .

We obtain the function as follows: First, write down the reverse Chern polynomials of E, F and E \bigotimes F with X, Y and T as indeterminates:

$$P(X) = X^{e} + c_{1}(E)X^{e-1} + \dots + c_{e}(E)$$

$$Q(Y) = Y^{f} + c_{1}(F)Y^{f-1} + \dots + c_{f}(F)$$

$$R(T) = T^{ef} + c_{1}(F \otimes F)T^{ef-1} + \dots + c_{ef}(E \otimes F) .$$

Now regard the coefficients of P and Q as transcendentals. Then in some field extension of Q we have

$$P(X) = \prod_{i} (X - \ell_{i})$$
$$Q(Y) = \prod_{i} (Y - m_{i})$$
$$\mathbf{j}$$

This being so, it now follows from $[F\ell]$ or any other standard source on Chern classes that

$$R(T) = \prod_{i,j} (T - (\ell_i + m_j)) .$$

For more details, see [Hm 6] .

the Chern polynomial of E

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the rank of E

 $TE_{-} = 1 + 0, (F)T + ... + 0, (E)T$

the Chern polynomial of F

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the rank of f; will return the Chern polynomial of Figs F in tarms

 $c_1(\Xi)$, \ldots $c_n(\Xi)$, $c_1(\Xi)$, \ldots $c_n(\Xi)$, $c_n(\Xi)$

We obtain the function as follows: First, write down the reverse Chern polynomials of E, F and E F with X, Y and T

 $Q(\mathbf{X}) = \mathbf{X}^{0} + o_{1}(\mathbf{X})\mathbf{X}^{0-1} + \dots + o_{n}(\mathbf{X})$ $Q(\mathbf{X}) = \mathbf{Y}^{0} + o_{1}(\mathbf{X})\mathbf{Y}^{0-1} + \dots + o_{n}(\mathbf{X})$ $\mathbf{R}(\mathbf{X}) = \mathbf{Y}^{0} + o_{1}(\mathbf{X})\mathbf{Y}^{0-1} + \dots + o_{n}(\mathbf{X})$

Now regard the coefficients of P and Q as transcendentals. Then in some field extension of Q we have

$$P(X) = \prod_{i=1}^{n} (X - i_{i})$$

This being so, it now follows from [F4] or any other standard source on -Chern classes that

$$R(T) = \prod_{i=1}^{m} (T - (I_i + m_j)) .$$

For more details, see [Hm 6] .

Hence it is natural to introduce the relation

T = X + Y ,

and we get

$$Q(T - X) = (-1)^{f} \prod_{j} (X - (T - m_{j}))$$

Recall, [VdWa] Vol I section 28, that the resultant of two polynomials

$$f(X) = a_0 X^n + a_1 X^{n-1} + \dots + a_n$$

$$g(X) = b_0 X^m + b_1 X^{m-1} + \dots + b_m$$

where a_0 and b_0 are non-zero, is equal to

$$r = a_0^m b_0^n \prod_{i,j} (x_j - y_i) ;$$

 x_1, \ldots, x_n and y_1, \ldots, y_m being all the roots of

f(X) = g(X) = 0

in some splitting field. Thus letting

$$f(X) = Q(T - X)$$
, $g(X) = P(X)$,

we get that

$$r = (-1)^{f} \prod_{i,j} (T - m_{j} - l_{i}),$$

so that up to a sign, the resultant of Q(T - X) and P(X) with respect to X is equal to R(T).

The function RESULTANT in MACSYMA may use different algorithms, see [MAC] p. 118. Normally the usual determinant is not computed directly. This may in some cases yield a sign different from what one expects. Rather than to keep track of this, it is better to adjust the sign in the end, which is easy because any Chern polynomial has 1 as its constant term. Hence it is natural to introduce the relation

· X + X = T

dog av bns

$$Q(T - X) = (-1)^T \prod (X - (T - m_j))$$
.

Recall, (Vows) Vol I section 28, that the resultant of two

$$x^{a} + \dots + \frac{1-\alpha}{2}x^{a} + \frac{\alpha}{2}x^{a} = (x)^{a}$$

 $x^{a} + \dots + \frac{1-\alpha}{2}x^{a} + \frac{\alpha}{2}x^{a} = (x)^{a}$

where and bo are non-zero, is equal to

x, ... xn and y, yn being all the roots of

$$f(x) = g(x) = 0$$

in some splitting field. Thus latting

$$f(X) = Q(T - X)$$
, $g(X) = P(X)$,

we get that

$$(1^{1} - 1^{2} - 1^{2}) = 1^{2}$$

so that up to a sign, the resultant of Q(T - X) and P(X) with respect to X is equal to R(T).

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TENSOR (E, M, F, N): =BLOCK ((L, P, Q, R, S, U, V), L:M*N, P:X^M*SUBST(1/X, T, E), Q:(T-X)^N*SUBST(1/(T-X), T, F), R:RESULTANT(P, Q, X), S:T^L*SUBST(1/T, T, R), S:RATEXPAND(S), U:S*COEFF(S, T, 0), V:SUM(COEFF(U, T, I)*T^I, I, 0, L), V);

In order to demonstrate this, we wish to generate Chern polynomials of two Modules, denoted by A and B. We do this by the array-defined function

A[J]:=1+SUM(CONCAT(C,I,A)*T^I,I,1,J)\$

and a similar definition	involving B. The result is as follows:
(C12) A[1]; (D12)	C1A T + 1
(C13) A[2];	C18 C20 + C1A C28 + C2A + C18. C2A
(D13)	$\begin{array}{c}2\\C2A T + C1A T + 1\end{array}$
(C14) A[3];	
(D14)	$\begin{array}{cccc} 3 & 2 \\ C3A T + C2A T + C1A T + 1 \end{array}$
and for B,	
(C15) B[1]; (D15)	C1B T + 1

(C16)	B[2];	ted for processing "by h
(D16)		C2B T + C1B T + 1
(C17)	B[3];	to objain Supetions whi
(D17)		C3B T + C2B T + C1B T + 1

If for instance F is of rank 1 , then of course we have a well known formula for the Chern classes of $E \bigotimes F$, namely

 $c_{k}(E \otimes F) = \sum_{i=0}^{k} {e-k+1 \choose i} c_{k-i}(E) c_{i}(F)^{i}$.

- 4 -

Our function TEMBON is written as follows:

```
TENSOR (E. M.F. M. I. EL OCKTOL, F.O. A.S.U. VI.I. AMA,

P. X^MASUBST (I/X. 1.E). 0: (T-X) *MSUBST (I/(T-X). 1.F),

B: BESULTAM (P.O.X). SI I'L #SUBST (I/T.T.R). B: RATERPAND (S).

U: S=COCCFF (S. T. B). VI SUBJCICEFF (U.T. 1) #T'J. F. D.L). VII
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In order to demonstrate this, we wish to generate Chann polynomials of two Modulas, denoted by A and B. We to this by the array-defined fun-tion

If for instance F is of rank 1, then of course we have a well nown formula for the Cheiro classes of I CO F ; hanely

c_k(z⊗z) = ∑((+++))c_{k-4}(z)c₃(z))²

. .

Moreover, for e = f = 2 the polynomial is rather simple, and may be familiar. We get:

(C18) TENSOR (A (1), 1, B (1), 1); (C1B + C1A) T + 1 (C19) TENSOR (A (1), 1, B (2), 2); (D19) (C2B + C1A C1B + C1A) T + (C1B + 2 C1A) T + 1 (C20) TENSOR (A (1), 1, B (3), 3); (D20) (C3B + C1A C2B + C1A C1B + C1A) T + (C2B + 2 C1A C1B + 3 C1A) T + (C1B + 3 C1A) T + 1 (C21) TENSOR (A (2), 2, B (2), 2); (D21) (C2B - 2 C2A C2B + C1A C1B C2B + C1A C2B + 2 C1B C2A + C1B C2A + C1A C1B C2A) T + (2 C1B C2B + 2 C1A C2B + 2 C1B C2A + 2 C1A C2A + C1A C1B² + C1A² C1B) T + (2 C2B + 2 C2A + C1B + 3 C1A C1B + C1A) T + (2 C1B + 2 C1A) T + 1

However, already for e = f = 3 the expressions become quite formidable, and completely unsuited for processing "by hand". See the appendix, section 1.

We may use a similar method to obtain functions which return the Chern polynomials of any exterior or symmetric power as well. This will not be needed here, but to illustrate we shall give a function

EXTERIOR 2 (arg₁, arg₂),

which when given the arguments

 $\arg_1 = 1 + c_1(E)T + \ldots + c_e(E)T^e$,

the Chern polynomial of E in Which T must be the indeterminate, and

Moreover, for s = f = ? the polynomial is rather simple, and may be familiar. We get:

 C18)
 TENSOR(ALL), I, B(L), I);

 C19)
 TENSOR(ALL), I, B(L), Z);

 C19)
 TENSOR(ALL), I, B(L), Z);

 C19)
 TENSOR(ALL), I, B(L), Z);

 C20)
 TENSOR(ALL), Z);

 C20)
 TENSOR(ALL), Z);

 C20)
 TENSOR(ALL), Z);

 C20)
 TENSOR(ALL), Z);

C21) TENSOR (A [2], 2, 8 [2], 2); D21) (C28 - 2 C2A C28 + C1A C19 C28 + C1A C28 + C2A + C18 C2A + C1A C18 C2A) T + '12 C18 C28 + 2 C1A C29 + 2 C18 C2A + 2 C1A C2A + C1A C18 + C1A C18) T + '12 C28 + 2 C2A + C18 + 3 C1A C18 + C1A + T + C1A C18) T + '12 C19 C28 + 2 C2A + C18 + 3 C1A C18 + C1A + T + C1A C18 + 2 C1A) T + '

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We may use a similar method to obtain functions which return the Chern polynomials of any exterior or symmetric power as well. This will not be needed here, but to illustrate we shall give a function

when given the secondents

T(E) = 1 + c1(E) + 1 + ce(E)

the Gnern polynomial of E in which T must be the indeterminate,

$$\arg_2 = e$$
,

the rank of E which must be ≥ 2 ; will return the Chern polynomial of $\Lambda^2 E$. It is clear from what follows how to generalize this to a function

which when given a third argument

$$\arg_{z} = r$$
,

which must be $\leq e$; will return the Chern polynomial of $\Lambda^r E$. The function is:

> EXTERIOR2(E, M):=BLOCK([L, B, P, O, R, S, U, V], L:M^2, B:M*(M-1)/2, P:X^M*SUBST(1/X, T, E), Q:SUBST(T-X, X, P), P:RATEXPAND(P), Q:RATEXPAND(Q), R:RESULTANT(P, Q, X), S:R/(2^M*SUBST(1/2*T, X, P)), S:RATEXPAND(S), U:FACTOR(S), V:SQRT(U), V:PART(V, 1), V:T^(M*(M-1)/2)*SUBST(1/T, T, V), V:RATSIMP(V), V:SUM(COEFF(V, T, I)*T^I, I, 0, M*(M-1)/2), V);

To see why this yields the correct result, denote the reverse Chern polynomials of E, E \bigotimes E and Λ^2 E by, respectively P(T), Q(T) and R(T). As above, writing

$$P(T) = \prod_{i} (T - \ell_{i}) ,$$

we have

$$Q(T) = \prod_{i_1, i_2} (T - (\ell_{i_1} + \ell_{i_2}));$$

and

$$R(T) = \prod_{i_1 < i_2} (T - (\ell_{i_1} + \ell_{i_2}));$$

the latter being the standard way in which the Chern classes of exterior powers are determined. Using this, we get

1 9 = 0318

the rank of E which must be $E \in I$, will return the Chern polynomial of A^2E . It is clear from what follows how to generalize this to a function

(para cars , are) MOIRETXE

which when given a third argument

· 2 = "318

which must be s e; will return the Chern polynomial of A'B

XTERIOR2 (E, M): -BLOCK (LL, B, P, O.R.S. U, VI, L'IMP2, B, M& H-L) / 2.
 PEXMESSION (F, N. T. EF, OLSOST (T-X, X. P), P. RATEXPAND (P), OLSOST (T-X, X. P), P. RATEXPAND (P), SERVER (P)

To see why this yields the correct result, denote the reverse thern polynomials of E . E \otimes E and A^2 E by, respectively P(T), Q(T)

we have

 $(T) = \prod_{l \neq 1} (T - (l_{11} + l_{12})) t$

bris

 $u((_{24}u + _{44}u) - T)_{24} = (T)u$

the latter being the standard way in which the Ghern classes of exterior powers are determined. Using this, we get

$$Q(T) = \prod_{i} (T - 2\ell_{i}) \{ \prod_{i_{1} \leq i_{2}} (T - (\ell_{i_{1}} + \ell_{i_{2}})) \}^{2}$$

 $= 2^{e} P(\frac{1}{2}T) R(T)^{2}$.

This should explain everything in the function body above, except possibly the use of the function PART. This is necessary since SQRT applied to $R(T)^2$ will return ABS(R(T)). The PART-function picks out the expression R(T).

We obtain the following polynomials for $2 \le e \le 4$:

(C23) EXTERIOR\2(A[1].1): Part fell off end. (C24) EXTERIOR\2(A[2],2); (D24)C1AT + 1(C25) EXTERIOR\2(A[3],3); 3 2 2 (D25) - C3A T + C1A T + 2 C1A T + 1 (C26) EXTERIOR\2(A[4].4): You have run out of LIST space. Do you want more? Type ALL; NONE; a level-no. or the name of a space. ALL; 2 2 6 3 3 - C3A T + C1A T + 3 C1A T + 3 C1A T + 1 (D26)

As we see, the improper arguments given to the function on line (C 23) results in an error message. The function call on line (C 26), while returning a simple polynomial as the answer, clearly generates large intermediate expressions.

Unfortunately this is not an infrequent situation, which tends to limit the range of results obtainable by methods such as the ones developed in this paper.

$a(\tau) = \prod_{i=1}^{n} (\tau - 2i_{1}) \prod_{i=1}^{n} (\tau - (z_{i_{1}} + z_{i_{2}})))$

This should explain everything in the function body above, except possibly the use of the function PART. This is necessary since SQRT applied to $R(T)^2$ will return ASS(R(T)). The PART-function picks out the expression R(T).

We obtain the following polynomials for 2 % e * 4 ;

(C23) EXTERIOR\2(A111,1); Part feil off end. (C24) EXTERIOR\2(A121,2); (C25) EXTERIOR\2(A12),3); (C25) EXTERIOR\2(A13),3); (C25) EXTERIOR\2(A14),4); (C25) EXTERIOR\2(A14),4); (C26) EXTERIOR\2(A14),4); (C27) EXTERIOR\2(A14

(D28) - C3A T + C1A T + 3 C1A T + 3 C1A T + 1

As we see, the improper arguments given to the function on line (C 23) results in an error message. The function call on line (C 25), while returning a simple polynomial as the answer, clearly generates large intermediate appressions.

Unfortunately this is not an infrequent situation, which tends to limit the range of results obtainable by methods such as the ones developed in this paper.

§ 2. Chern- and Segre classes of Grassmanians.

Let $G_k(r, A) = G(r, A) = G$ be the Grassmanian which parametrizes the linear r-subspaces of $\mathbb{P}_k^A = \mathbb{P}^A$. Equivalently, G(r, A) is the scheme which represents the functor of r + 1 - quotients of $V = k^{A+1}$, so in particular it carries the universal quotient

$$V_{G} \longrightarrow Q \longrightarrow 0$$

where Q is locally free of rank r + 1.

Here (i.e. in this and the following section), we only summarize the basic formulae needed from the theory of Grassmanians in algebraic geometry, and the closely related theory of Schubert Calculus. For details, the reader is referred to [Hm 7] or the references given there, such as [Lk 1]. Last but not least, the recent paper of <u>R. Donagi</u> [Do] contains among other things an exellent account of some of the material we need here.

The reader not familiar with the material which follows, may consult [Hm 7] for a more extensive survey, with references to the literature. In using these, one should of course beware of the distinction between the projective and the affine notation. Thus the Grassmanian which we denote by G(r, A) would be denoted by G(r + 1, A + 1) in affine notation, used for instance in [Do].

First, we have the basic formula (2.1) $\Omega_G^1 = Q^V \otimes M$, M being defined by the exact sequence (2.2) $0 \longrightarrow M \longrightarrow V_G \longrightarrow Q \longrightarrow 0$ and where

$$Q^{\vee} = \underline{Hom}_{G}(Q, O_{X})$$

6 2. Chern- and Segre classes of Grassmanians.

Let $G_{\mathbf{x}}(\tau, \Lambda) = \mathbf{U}(\tau, \Lambda) = 0$ be the Grassmanian which parametrizes the linear r-subspaces of $\mathbf{F}_{\mathbf{x}}^{\Lambda} = \mathbf{F}^{\Lambda}$. Equivalently, $\mathbf{D}(\tau, \Lambda)$ is the scheme which represents the functor of $\tau + 1$ - quotients of $\nabla = \lambda^{1+1}$, so in particular it carries the universal quotient

where Q is locally free of rank T + 1

Here (1.e. in this and the following section), we only summarize the basic formulae needed from the theory of Grassmanians in algebraic geometry, and the closely related theory of Schubert Calculus. For details, the reader is referred to [im 7] or the references given there, such as [im 1]. Last but not least, the resent paper of <u>R. Donari</u> [Do] contains among other things an exellent account of some of the material we need here.

The reader not familiar with the material which follows, has consult (Hm 7) for a more arteneive survey, with references to the literature. In using these, one should of course beware of the distinction between the projective and the affine notation. Hous the Grassmanian which we denote by G(r, A) would be denoted by G(r + 1, A + 1) in affine notation, used for instance in [D0].

(2.1) no - 0 0 M.

M being defined by the exact sequence (2.2) $0 \rightarrow M \rightarrow V_{0} \rightarrow 0 \rightarrow 0$

. 8 .

is the dual of Q. Thus in particular G is non-singular and

$$\dim(G) = (r + 1)(A - r)$$
.

Moreover, as the sequence (2.2) is split we have the identification

$$G(r, A) = G(A - r - 1, A)$$

and it thus suffices to consider the cases

$$r \leq \left[\frac{A-1}{2}\right]$$
.

Since the case

$$G(0, A) = \mathbb{P}^{A}$$

is trivial as far as this investigation is concerned, we shall assume that

$$1 \leq r \leq \left[\frac{A-1}{2}\right]$$

Using the function TENSOR of the previous paragraph, together with the well known relation

$$c_t(Q) = c_{(-t)}(Q^{\vee})$$

where as always $c_t(E)$ denotes the Chern polynomial of E, it is now easy to write a function in MACSYMA which computes the Chern polynomial of G(r, A), i.e. the Chern polynomial of

$$T_{G} = \Omega_{G}^{1 \vee}$$

First, we introduce the function

CHERNPOLYBUNDLE(arg₁, arg₂, arg₃, arg₄),

which returns the Chern polynomial in the indeterminate given by \arg_4 , of a bundle (i.e. a locally free Module) denoted by \arg_1 of rank = \arg_2 on a variety of dimension = \arg_3 . We include also the function

SEGREPOLYBUNDLE(arg₁, arg₂, arg₃, arg₄),

which when given the same arguments as CHERNPOLYBUNDLE above, will return the Segre polynomial in terms of the Chern classes of the bundle. is the dual of Q. Thus in particular Q is non-singular and discrete Q.

Moreover, as the sequence (2,2) is split we have the identification a(r, s) = 0(A - r - 1, s)

and it thus suffices to consider the cases

is trivial as far as this investigation is concerned, we shall adount

$$I \in \mathbb{Z} \times \left[\frac{A}{2} \right]$$

Using the function TENSOR of the previous paragraph, together with the well known relation

where as always $d_{i}(E)$ denotes the Onern polynomial of E, it is now easy to write a function in MAGSYMA which computes the Chern polynomial of G(r, A), i.e. the Chern polynomial of

First, we introduce the function

CHERNFOLTBUNDLE(arg, srg, arg, arg,)

which returns the Ghern polynomial in the indeterminate given by \arg_{i_1} , of a bundle (i.e. a locally free Module) denoted by ang, of rank = \arg_2 on a variaty of dimension = \arg_2 . We include also the function

SEGREPOLYBUNDLE (arg, arg, arg, argg)

which when given the same arguments as GHERNPORTBUNDLE above, will return the Segre polynomial in terms of the Chern-classes of the bundle SEGREPOLYBUNDLE (Q, RANK, DIM, T):=TAYLOR(1/CHERNPOLYBUNDLE(Q, RANK, DIM, -T), T, Ø, DIM);

Furthermore, the function

CHERNPOLYMDUAL (arg, arg, arg,)

will return the Chern-polynomial with \arg_3 as the indeterminate of the bundle M^V on G(r, A), when given the arguments $\arg_1 = r$, $\arg_2 = A$. The result is expressed in terms of the Chern classes of Q. The function is:

CHERNPOLYMDUAL (R, A, T):=BLOCK ([RANK, DIM, P], RANK:R+1, DIM:A-R, P:SEGREPOLYBUNDLE (Q, RANK, DIM, T), EXPAND (P));

Using the above CHERNPOLYBUNDLE and CHERNPOLYMDUAL, together with TENSOR, we now write a function

POLYS(arg₁, arg₂),

which when given the arguments

 $\arg_1 = r, \arg_2 = A;$

will do the following: First, the Chern polynomial of G(r, A) is computed by the function CHERNPOLYG given below and assigned to the variable CHEPOrA. Next, the result is printed out as

"CHEPOrA =
$$c_+(G(r, A))$$
"

Finally the same is done for the Segre polynomial of G(r, A). The function is:

CHERNPOLYBUNDLE (O. RANK, DJ N. T) :-BLOCK (IP). Pr1+SUR(CONCAT(C, I, D) *T^I, I, I, HIN (RANK, DJ H) , PI

SECREPOLYBUNGLE (D. RANK, DIN, 11: - TAYLORI I ACHERIPOLYBUNGLE (D. RANK, DIN, - TI. T. B. DIN);

Parthermore, the Function

CHERNPOLYMOUAL (arg, arg, arg,

will return the Chern-polynomial with arg as the indeterminate of the bundle H^{\prime} on G(r, A), when given the arguments arg r_i , arg r A. The result is expressed in terms of the Chern classes of The function is:

CHERNPOLYHOUAL IR, A, T1: -BLOCK (IRAMK, DIM, P1, BANK, RAL, DIMLA-R, P. SEGREPOLYBUNDLE (D. RANK, DIM, T1; EXPAND/P3);

Using the above CHERNFOLYEUWIALS and CHERNFOLYMOUAL, together with

Forxs (arg, args) .

which when given the arguments

f w = Cars 'I = 'SIU

will do the following: First, the Chein polynomial of $O(\tau, A)$ is computed by the function CHEANFOLYC given below and assigned to the variable CHEFORA. Maxt, the result is printed out as "CHEFORA = $e_{e}(G(\tau, A))$ ".

Finally the same is done for the Segre polynomial of G(r, A). The function is:

The function CHERNPOLYG(\arg_1 , \arg_2 , \arg_3) which computes the Chernpolynomial of G(r, A) in terms of the Chern classes of Q, with \arg_3 as the indeterminate and where

$$\arg_1 = r, \arg_2 = A$$
,

is as follows:

CHERNPOLYG(R, A, T):=BLOCK([RQ,DIM,RM,F,G,H],RQ:R+1, DIM:(R+1)*(A-R),RM:A-R,F:CHERNPOLYBUNDLE(Q,RQ,DIM,T), G:CHERNPOLYMDUAL(R,A,T),H:TENSOR(F,RQ,G,RM),H);

Note that M^{\vee} is of rank A - r, so the polynomial $c_t(M^{\vee})$ is of degree A - r. This observation yields the set of relations (actually: a set of generators for the ideal of relations) among the Chern classes of Q: In fact, the inverse of the polynomial $c_t(Q^{\vee})$, where $c_1(Q)$, ..., $c_{r-1}(Q)$ are regarded as transcendentals for the moment, contains terms

 $\left\{\rho_{i}(c_{1}(Q), \ldots, c_{r+1}(Q))t^{i} \mid A - r < i \leq (r+1)(A - r)\right\}$ and these are all equal to zero, thus giving r(A - r) relations among the Chern classes of Q. It is a classical result that these are all the relations.

This observation enables us to write the function

RELATIONSOFCHERNCLASSES(arg1, arg2),

which when given the arguments

POLYSG (R:A) := BLOGRE MINT CH. 0.561 (MINT HV) A ATI : CH. CHI CHERRIP & CH. A. T) .CONCATI'CHERD, R. ATI : CH. PRINT ("CHERD". R. A. "-". CH) .0:1/CH. 56: TAV. CR (D. T. 4. 01M) BE: 1:580016XCAND (COEF (SE. T. 1) AT" 1.1.1,01M) CONCAT(SEPO. R. A) :: SE. PRINT ("SEPO", R. A. "-". SE) .01M1

The function (HERNFOLMG(arg₁, arg₂, arg₃) which computes the Cherngolynomial of $O(\tau, A)$ in terms of the Chern classes of Q, with arg₃ as the indeterminate and where arg₃ as the indeterminate and where

is sollows:

CHERNPOLYG (R, A, T) : - OL DOX (180, 010, R0, F, C, H), R0:R+L. 010: (R-1) = (A-R), R0: A-R, R: CEERIPOLY00002.ER. A0, 010, T+. 6: CHERNPOLYDOAL (R, A, T), H: TEREOR (F, PO.C. B1), H);

Note that M^{V} is of rank A - r, so the polynomial $c_{c}(M^{V})$ is of degrees A - r. This observation yields the set of relations (actually: a set of generators for the ideal of relations) among the Ghern classes of Q r. In fact, the inverse of the polynomial $c_{l}(Q^{V})$, where $c_{1}(Q)_{1}, \dots, c_{r-1}(Q)$ are regarded as transcendentals for the

 $\begin{cases} e_1(e_1(0), \dots, e_{n+1}(0))t^n \mid A - T \leq 1 \leq (T + 1)(A - T) \\ \text{and these are all equal to zero, thus giving <math>T(A - T)$ relations anong the Gnern classes of Q. It is a classical result that these are all the relations:

This observation enables us to write the function RELATIONSOFCHERNCLASSES(arg., arg.) .

which when given the argument.

 $\arg_1 = r$, $\arg_2 = A$;

will return a print out as follows

" $\rho_i(c_1(Q), \ldots, c_{r+1}(Q)) = 0$ "

for i = A - r + 1, ..., (r + 1)(A - r). We are not using this in the sequel, otherwise it would of course be better to let the function return a <u>list</u> of the relations, which could then be further manipulated on by other functions. We have:

> RELATIONSOFCHERNCLASSESG(R,A):=BLOCK([RANK,RANKM,P,RE],RANK:R+1, RANKM:A-R,DIM:RANK*RANKM, P:SEGREPOLYBUNDLE(0,RANK,DIM,T), FOR I FROM RANKM+1 THRU DIM DO (RE:EXPAND(COEFF(P,T,I)), PRINT(RE," = 0 ")));

The following function generates the Chern and the Segre polynomials, as well as the relations of the Chern classes of Q for G(r, A) for all

 $1 \leq r \leq \left[\frac{A-1}{2}\right]$:

G(A):=BLOCK([E],E:ENTIER((A-1)/2), FOR I THRU E DO (POLYSG(I,A),RELATIONSOFCHERNCLASSESG(I,A)));

For A = 3 we get the following:

are, = r , aregin a s

will return a print out as follows

 $1 = A - r + 1, \dots, (r + 1)(A - r)$. We are not using

return a list of the relations, which could then be further manipulated on by other functions. We have:

RELATIONSOF CHERNOLASSESS (R.AI: -8LOCK (RANK, RANNE, P. REI, RANK (Rel RANKTI: A-R.OIT: RANK REPORT P: SECREPOLYBUNOLE (C. RANK, DIM.T). FOR 1 FROM RANKTH: THEU DIM DO (RE: EXTAND (COEFF (P. T.11). PRIMITE: - - 8 711).

The following function generates the Chern and the Segre polynomials, as well as the relations of the Chern cleases of Q for G(r. A) for all

 $1 \le r \le \left[\frac{A-1}{2}\right]$:

G (A1 = BLOCK (16) . EFENTIER (IA - 1) A21 . FOR 1 THRU E DO (POLYSG (1, A) . RELATIONSOFCHERINGLASSERG (1, A111 ;

For A = 3 we get the following:

(C5) 6(3); CHEPO 1 3 - 14 620 - 4 C10 620 + 3 C10 1 1 + 6 G10 T + 7 C10 T CHEPO 1 3 - 14 620 - 4 C10 620 + 3 C10 1 1 + 6 G10 T + 7 C10 T

1 + 1 012 + - ⁵ 1.⁵ b12 + +

620 - 3 CIO C20 + CIO - 8

A full list of results for $3 \le A \le 6$ is given in the appendix, § 2.

There is a somewhat different method for computing the Chernand Segre classes of Grassmanians.

This utilizes the exact sequence

$$(2.3) \qquad 0 \longrightarrow \Omega^{1}_{G} \longrightarrow (Q^{V})^{A+1} \longrightarrow Q \otimes Q^{V} \longrightarrow 0$$

obtained from (2.2) by tensoring with Q^V . It yields the formulae

(2.4)
$$c_t(Gr, A) = (1 + c_1 t + \dots c_{r+1} t^{r+1})^{A+1} (c_t(Q \otimes Q^V))$$

(2.5) $s_t(G(r, A) = c_t(Q \otimes Q^V)/(1 + c_1t + ... + c_{r+1}t^{r+1})^{A+1}$

Using the function TENSOR, it is of course clear how to write functions which compute the polynomials according to the formulae above. It is best to do so by an array-defined function,

RE[arg] ,

which when given the argument $r = \arg$ will return $c_t(Q \otimes Q^{\vee})$ for Q of rank r + 1. Thus once RE is computed for a given value of r, it is stored and may be used again when needed for a different value of A. We delete the details. One then uses the function

> SEGREPOLYGRASS(R,A):=BLOCK([DIM,S],DIM:(R+1)*(A-R), S:RE[R]/(1+SUM(CONCAT(C,I,Q)*T^1,I,1,R+1))^(A+1), S:TAYLOR(S,T,0,DIM), S:1+SUM(EXPAND(COEFF(S,T,I))*T^I,I,1,DIM),S);

We obtain the following results, via a function

A full list of results for 3 5 A 5 5 is given in the appendix,

There is a somewhat different method for computing the Chernand Segre classes of Grassmanisma.

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obtained from (2.2) by tensoring with 9 . It yields the formulae

$$(2.4) \qquad c_{g}(a_{2}, A)) = (1 + c_{1}t + \dots + c_{r+1}t^{r+1})^{A+1}(c_{g}(a_{2}a^{2}))$$

Using the function TENSOR, it is of course clear how to write functions which compute the polynomials according to the formulae above. It is best to do so by an array-defined function,

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which when given the argument r = arg will return $c_{i}(0, 0, 0)$ for Q of rank r + 1. Thus once RE is computed for a given value of r, it is stored and may be used again when needed for a different value of A. We delete the details. One then uses the function

> BEGREPOLYGRASS (R. A) :-RECKX (1010, St. DHb (R+1) * (A-R), B: RE (R) / (1+SURVCONCATIC, F. 0) * (7: [, 1, 1, R+1)) * (A+1), S: TAYLOR (S. 1, 3, 01M), S: 1+SURVEXPAND (COEFF (S. 7, 11) * (7: [, 1, 1, 21M), S);

> > de obtain the following results, via a function

GG(arg)

which when given the argument

arg = A

will print out the new Chern and Segre polynomials in an analogous way to the function G(arg).

```
(CG) GG(3);

    2
    2
    4
    4

    NEWCHEPO 1 3 = (6 C2Q - 16 C1Q C2Q + 8 C1Q) T

                                                                                                                                                                                    \begin{array}{c} 3 \\ + (8 C10 \\ 2 \\ 2 \\ \end{array} \begin{array}{c} 3 \\ - 4 C10 \\ 2 \\ \end{array} \begin{array}{c} 3 \\ T \\ + 7 \\ \end{array} \begin{array}{c} 2 \\ - 7 \\ \end{array} \begin{array}{c} 3 \\ - 7 \\ - 7 \\ \end{array} \begin{array}{c} 2 \\ \end{array} \begin{array}{c} 2 \\ - 7 \\ \end{array} \begin{array}{c} 2 \\ \end{array} \begin{array}{c} 2 \\ \end{array} \begin{array}{c} 2 \\ - 7 \\ \end{array} \begin{array}{c} 2 \\ \end{array} \begin{array}{c} 2 \\ - 7 \\ \end{array} \begin{array}{c} 2 \\ \end{array} \end{array} \begin{array}{c} 2 \\ \end{array} \begin{array}{c} 2 \\ \end{array} \begin{array}{c} 2 \\ \end{array} \begin{array}{c} 2 \\ \end{array} \end{array} \begin{array}{c} 2 \\ \end{array} \begin{array}{c} 2 \\ \end{array} \end{array} \begin{array}{c} 2 \\ \end{array} \begin{array}{c} 2 \\ \end{array} \end{array} \end{array} \begin{array}{c} 2 \\ \end{array} \end{array} \begin{array}{c} 2 \\ \end{array} \end{array} \begin{array}{c} 2 \\ \end{array} \end{array} \end{array} \begin{array}{c} 2 \\ \end{array} \end{array} \begin{array}{c} 2 \\ \end{array} \end{array} \end{array} \end{array}
 NEWSEPO 1 3 = (-6 C20 - 16 C10 C20 + 25 C10) T
                                                                                                                                                                             3 3 2 2
+ (4 C10 C20 - 16 C10 ) T + 9 C10 T - 4 C10 T + 1
    (D6)
                                                                                                                                                                                                                                                                                                                                    DONE
    (C7) GG(4):
 NEWCHEPO 1 4 = (-14 C20 + 88 C10 C20 - 72 C10 C20 + 16 C10) T
         2 3 5 5
+ (30 C10 C20 - 40 C10 C20 + 16 C10 ) T
         + (6 C2Q - 13 C1Q C2Q + 16 C1Q) T + 15 C1Q T
 + (C20 + 11 C10) T + 5 C10 T + 1
3 2
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                              6 6
NEWSEPO 1 4 = (25 C20 - 15 C10 C20 - 245 C10 C20 + 148 C10 ) T
         + (15 C10 C20 + 110 C10 C20 - 91 C10 ) T
         + (- 5 C2Q - 40 C1Q C2Q + 55 C1Q ) T + (10 C1Q C2Q - 30 C1Q ) T
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                         3
      2 2
+ (14 C10 - C20) T - 5 C10 T + 1
                                                                                                                                                                                                                                                                                                                                    DONF
```

A full list of results for $3 \le A \le 6$ is given in the appendix, § 2.

We see that the Chern and Segre polynomials returned by the two methods are not identical. This is of course due to the relations among the Chern classes of Q.

As we shall see in following paragraphs, it will be essential to find the simplest possible expressions for the Cern and Segre classes, at least when r and A are large.

The above data seem to indicate that the first method yields a simpler result than the second. However, for small values of r the second method is the best. In fact, once RE[R] is given, we can compute the polynomials for any A using less sophisticated systems than MACSYMA, and even obtain explicit if messy formulae in some cases. We return to this in section 5.

There we will treat the case r = 1, i.e. G(1, N)'s in this manner, and see how the computations we are interested in, and where we use the Chern and Segre polynomials, reduce to the evaluation of straight forward combinatorical identities, which present no difficulties from a computational point of view.

Moreover, there should be a good possibility that the case r = 2 can be treated analogously, as we shall indicate. But for r = 3 there is no hope of obtaining similar combinatorical expressions, even though the method might still be usefull from a numerical point of view. This becomes clear from the expressions for RE[1], RE[2] and RE[3] which are listed below. For R = 4 the computation became too large for the system, in that the intermediate expressions filled up all available space.

Thus the first method for computing the Chern and Segre polynomials

- 15 -

A fall list of results for 3 = 4 = 6 is given in the appendix.

We see that the Ghern and Segre polynomials returned by the two methods are not identical. This is of course due to the relations among the Chern classes of Q.

As we shall see in following paragraphs, it will be essential to find the simplest possible expressions for the Cern and Sogre classes,

The above data seem to indicate that the first method yields a simpler result than the second. However, for small values of f the second method is the best. In fact, once RE[R] is given, we can compute the polynomials for any A using less sophisticated systems than MACSYMA, and even obtain explicit if meany formulae in some cases. We return to this in section 5.

There we will treat the case real, i.e. G(T, M)'s in this manner, and see how the computations we are interested in, and where we use the Chern and Sagre polynomiels, reduce to the evaluation of straight forward combinatorical identifies, which present no difficulties from a computational point of view.

Moreover, there should be a good possibility that the case T = 2can be treated enalogously, as we shall indicate. But for T = 3there is no hope of obtaining similar combinatorical expressions, even though the method might still be useful from a numerical point of view. This becomes thear from the supressions for RE[1], RE[2] and RE[3] which are listed below. For R = 4 the computation become too large for the system, in that the intermediate expressions filled up all available space.

Thus the first method for desputing the Ohern and Segre polynomials

15

would yield new information for the first time for G(4,9).

But at this level the method appears to be impracticable on MACSYMA. I say <u>appear</u>, because I have not tried to have this done as a background job with disk use.

My present opinion is that it would be better to have the functions implemented directly in a lower level language.

(C11) RE[1]: 2 2 (4 C2Q - C1Q) T + 1(D11) (C12) RE[2]; 2 3 (D12) (27 C30 - 18 C10 C20 C30 + 4 C10 C30 + 4 C20 - C10 C20) T 2 2 4 4 2 2 + (9 C2Q - 6 C1Q C2Q + C1Q) T + (6 C2Q - 2 C1Q) T + 1 (C13) RE[3]; You have run out of LIST space. Do you want more? Type ALL; NONE; a level-no. or the name of a space. ALL: 2 2 (D13) (256 C40 - 192 C10 C30 C40 - 128 C20 C40 + 144 C10 C20 C40 2 2 - 27 C1Q C4Q + 144 C2Q C3Q C4Q - 6 C1Q C3Q C4Q - 80 C10 C20 C30 C40 + 18 C10 C20 C30 C40 + 16 C20 C40 2 3 4 3 3 - 4 C10 C20 C40 - 27 C30 + 18 C10 C20 C30 - 4 C10 C30 2 2. 12 - 4 C2Q C3Q + C1Q C2Q C3Q) T 2 . + (- 192 C20 C40 + 72 C10 C40 + 216 C30 C40 - 120 C10 C20 C30 C40 3 3 2 2 3 + 18 C10 C30 C40 + 32 C20 C40 - 6 C10 C20 C40 - 54 C10 C30 4 .. 2 2 2 2 4 2 3 + 18 C20 C30 + 42 C10 C20 C30 - 9 C10 C30 - 26 C10 C20 C30 3 2 5 2. 4 10 + 6 C10 C20 C30 + 4 C20 - C10 C20) T

- 16 -

would yield new information for the first time for 0(4,9). But at this level the method appears to be impracticable on MACHYMA. I say <u>appear</u>, because I have not tried to have this done as a background tob with tisk use.

My present opinion is that it would be batter to have the Functions implemented directly in a lower level language.

+ $(-112 \ C40^{2} + 56 \ C10 \ C30 \ C40 + 24 \ C20^{2} \ C40 - 32 \ C10^{2} \ C20 \ C40$ + $6 \ C10^{4} \ C40 + 48 \ C20 \ C30^{2} - 25 \ C10^{2} \ C30^{2} - 54 \ C10 \ C20^{2} \ C30$ + $38 \ C10^{3} \ C20 \ C30 - 6 \ C10^{5} \ C30 + 17 \ C20^{4} - 12 \ C10^{2} \ C20^{3}$ + $2 \ C10^{4} \ C20^{2} \ T^{8} + (16 \ C20 \ C40 - 6 \ C10^{2} \ C40 + 26 \ C30^{2}$ - $38 \ C10 \ C20 \ C30 + 8 \ C10^{3} \ C30 + 28 \ C20^{3} - 24 \ C10^{2} \ C20^{2} + 8 \ C10^{4} \ C20$

6 6 - C1Q) T + (8 C4Q - 2 C1Q C3Q + 22 C2Q - 16 C1Q C2Q + 3 C1Q) T

2 2 + (8 C2Q - 3 C1Q) T + 1

(C14) RE[4];

You have run out of LIST space. Do you want more? Type ALL; NONE; a level-no. or the name of a space. ALL;

You have run out of LIST space. Do you want more? Type ALL; NONE; a level-no. or the name of a space. ALL;

CORE capacity exceeded (while requesting FIXNUM space)

33037 msec. so far

0² + 56 CTO COO CAO + 24 CZO CAO - 32

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§ 3. A basis for some Schubert Calculus on a computer.

Here we take the term Schubert Calculus to mean the formal computation with "Schubert Cycles" $\Omega(a_1, \ldots, a_q)$ in A(G(r, A)), where q = A - r.

Different notations are in use, for the ones utilized here we refer to [Hm 7]. Also, there the reader will find references to some of the literature, of which we rely particularly on the fundamental and classical source [HP].

We shall make no attempt here to pursue the Schubert Calculus on a computer for its own sake, even though such a project would certainly be a very interesting one. Rather, we develope the minimum regired for our present purpose.

Indeed, recall the Plücker-embedding

$$p : G(r, A) \hookrightarrow \mathbb{P}(\Lambda^{r+1} \vee) = \mathbb{P}^{N}$$

where

$$N = \begin{pmatrix} A + 1 \\ r + 1 \end{pmatrix} - 1$$

Then we have that

$$O_{G(r,A)}(1) = p*(O_{IP}(1)) = \Lambda^{r+1}Q$$
.

Hence in particular if D is a very ample divisor giving the embedding p , then

$$[D] = c_1(Q)$$
.

In the next sections we shall see how this makes it possible to describe properties of embeddings and duality for Grassmanians, as

- 18 -

S 3 A besis for some Schubert Calculus on a computer.

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Indeed, recall the Flücker-embedding

$$\mathbf{E} : \mathbf{G}(\mathbf{r}, \mathbf{A}) \xrightarrow{q_{1}} \mathbf{E}(\mathbf{A}^{r+1} \vee) = \mathbf{E}^{\mathbf{H}}$$

where

$$I = \begin{pmatrix} 1 & + & 4 \\ 1 & + & 1 \end{pmatrix} = I$$

Then we have that

$$O_{G(T,A)}(1) = p^{+}(0_{\underline{M}}(1)) = A^{T+1}q.$$

Hence in particular if D is a vary ample divisor giving the embedding p . then

In the next sections we shall see now this makes it possible to describe properties of embeddings and duality for Grassmanians, as well as for a large class of varieties "generated" by Grassmanians in a sense made precise later, in terms of the Chern numbers of the universal quotient bundles Q: That is to say, in terms of the degrees of the elements

$$c_1(Q)^{i_1} \dots c_{r+1}(Q)^{i_{r+1}} \in A^{\dim}(G(r, A))$$

where $i_1 + \ldots + i_{r+1} = \dim = (r + 1)(A - r)$. The degree map is in this case an isomorphism

$$A^{\dim}(G(r, A)) \xrightarrow{\deg} \mathbb{Z}$$

so we may refer to the momomials above as the Chern numbers of Q .

Hence we need an algorithm to compute the Chern numbers of Q for any G(r, A). For this we proceed as follows: It is possible to convert any polynomial F in $c_1(Q), \ldots, c_{r+1}(Q)$ with integral coefficients into a linear combination in the elements $\Omega(a_1, \ldots, a_q) \in A(G(r, A)), q = A - r$ and $1 \le a_1 < a_2 < \ldots < a_q \le A + 1$. This is done by repeated application of <u>Pieri's Formula</u>, which asserts the following:

(3.1)
$$\Omega(a_1, \ldots, a_q)c_h(Q) = \sum \Omega(b_1, \ldots, b_q)$$

where the sum is extended over all indicies for which

$$\sum_{j=1}^{q} b_j = \sum_{j=1}^{q} a_j - b_j$$

and

$$1 \leq b_1 \leq a_1 < b_2 \leq a_2 < \ldots < b_q \leq a_q \leq A + 1$$

Since

well as for a large class of variables "generated" by Gracemanians in a sense made precise later, in terms of the Gnern numbers of the universal quotient bundles Q: That is to say, is terms of the degrees of the elements

where $1_1 + \ldots + 1_{r+1} = dim = (r + 1)(A - r)$. The degree map is in this case an isomorphism

so we may refer to the momonfais above as the <u>Chern numbers of 9</u>. Hence we need an algorithm to compute the Chern numbers of 9 for any G(r, A). For this we proceed as follows: It is possible to convert any polynomial 7 in $c_1(Q), \dots, c_{r+1}(Q)$ with integral coefficients into a linear combination in the elements $\Omega(a_1, \dots, a_Q) \in \Lambda(G(r, A)), q - A - r and 1 \le a_1 \le a_2 \le \dots \le a_Q \le A + 1$ This is done by repeated application of <u>Fisti's Formula</u>; which accepts the following:

$$(a_1, \dots, a_n) = \sum_{i=1}^{n} a_{i} (a_1, \dots, a_n) = \sum_{i=1}^{n} a_i (a_1, \dots, a_n)$$

where the sum is extended over all indicies for which



bris

1 5 b, 5 a, c b, 5 a, c ... c b, s a, s a s

Since

 $1 = \Omega(r + 2, ..., A + 1)$

in A(G(r, A)) we obtain the algorithm in question as follows: First multiply F by $\Omega(r + 2, ..., A + 1)$, then perform the substitution (3.1) in F repeatedly until F is free of $c_{h}(Q)$, for h = 1, ..., r + 1.

Since there is only one
$$\Omega(a_1, \ldots, a_q) \in A^{\dim}(G(r, A))$$
, namely,
 $a_1 = 1, a_2 = 2, \ldots, a_q = A - r$

and this element is easily seen to have degree 1, the case when F is homogeneous of (weighed) degree dim will yield an integral multiplum of $\Omega(1, \ldots, A - r)$, and the numerical factor is the sought degree.

In principle this will solve our problem. However, it may be better to proceed in a slightly different way. This is due to the fact that the above procedure, when implemented in MACSYMA, tends to be quite time-consuming. What we can do, is first of all to observe that ______r

$$deg(c_{1}(Q)^{i}) = deg(G(r, A)) = \frac{\prod^{i}(i!)((r + 1)(A - r))!}{\prod^{A}(i!)}$$
$$= \frac{\prod^{A}(i!)}{1=A-r}$$

for all i ≤ dim . Here the degree is with respect to the Plücker embedding, [Hm 7]. Moreover, with the same interpretation of "deg", we have

$$deg(c_{1}(Q)^{i_{1}}c_{2}(Q)^{i_{2}} \dots c_{r+1}(Q)^{i_{r+1}}) = deg(c_{2}(Q)^{i_{2}} \dots c_{r+1}(Q)^{i_{r+1}})$$

and since there are formulae for computing the degrees of $\Omega(a_1, \ldots a_q)$

in A(G(r, A)) we obtain the algorithm in question as follows: First multiply F by R(r + 2, ..., A + 1); then perform the substitution (3.1) in F repeatedly until F is free of $e_1(R)$; for h = 1, ..., r + 1.

ince there is only one
$$\Omega(a_1, \ldots, a_q) \in A^{\alpha_1 m}(\Theta(x, A))$$
, namely,
 $a_1 = 1, a_2 = 2, \ldots, a_q = A - 1$

and this element is easily seen to have degree 1. the case when F is homogeneous of (weighed) degree dim will yield an incegral multiplum of $\Omega(1, \ldots, A - r)$, and the numerical factor is the rought degree.

better to proceed in a slightly different way. This is due to the fact that the above procedure, when implemented in MACEINA, tends to be quite time-consuming. What we can do, is first of all to observe

$$deg(o_1(Q)^{\frac{1}{2}}) = deg(d(r, A)) = \frac{1-1}{1-1}$$

$$\prod_{i=1}^{A} (a_i)$$

for all is dim . Here the degree is with respect to the Plücker embedding, [Hm 7]. Moreover, with the same interpretation of "deg", we have

and since there are formulae for computing the degrace of G(4, ... 9,)

in general, all we need to do is to reduce monomials of the form

$$c_{2}(Q)^{i_{2}} \dots c_{r+1}(Q)^{i_{r+1}}$$

to a linear combination in the $\Omega\,{}^{\prime}s$. We return to this below.

We now give the function which transforms a polynomial in the $c_i(Q)$'s to a linear combination in the Ω 's.

First, the function DOMAIN(\arg_1 , \arg_2) generates the list of all lists of indicies $[a_1, \ldots, a_q]$ where $q = \arg_1, \Sigma a_1 = \arg_2$ and $1 \le a_1 < \ldots < a_q$:

```
DOMAIN(M,N):=BLOCK([S,R,L],L:[],

IF N >= M*(M+1)/2

THEN (IF EQUAL(M,1) THEN L:[[N]]

ELSE FOR I THRU N DO

FOR S IN DOMAIN(M-1,N-I) DO

(IF S[M-1] < I

THEN (R:ENDCONS(I,S),

L:CONS(R,L)))),L);
```

The function

will, when given the arguments

$$arg_1 = r$$
, $arg_2 = A$, $arg_3 = N$;

return the list of all lists of indicies

$$[a_1, \ldots, a_{A-r}]$$

where

$$1 \leq a_1 < \ldots < a_{N_r} \leq A + 1$$
,

and

$$\sum a_i = N :$$

and the resource and to do is to reduce monomials of the form

to a linear combination in the $\Omega's$. We return to this balow. We now give the function which transforms a polynomial in the $o_1(Q)$'s to a linear combination in the $\Omega's$. First, the function DOMAIN(erg₁, srg₂) generates the list of all lists of indicies $[a_1, \dots, a_q]$ where $q = arg_1$. Est = arg₂ and be a. $< \dots < a_n$:

```
DOMAIN(H, N):-BLOCK((S.R.L)LL)

IF N >- D*(H+1)/2

THEN (IF EQUAL (D.1) THENLLIIM)

ELSE FOR 1 THEN N DO

FOR 5 IN DOMAINMH-1, N-11 DO

THEN (R.ENDCOMS(1.5).

THEN (R.ENDCOMS(1.5).LL)
```

The function

DOMAING (arg, arg, arg,

will, when given the arguments

arg, = r, arg₂ = A, arg₃ = n

return the list of all lists of indicies

· lyng · · · · · · [8].

where

1 ± A & S Barry & A + 1

bras

DOMAING(R, A, N):=BLOCK([S, RR,L,M],L:[],M:A-R, IF N >= M*(M+1)/2 THEN (IF EQUAL(M,1) THEN (IF N <= A+1 THEN L:[[N]]) ELSE FOR I THRU MIN(N,A+1) DO FOR S IN DOMAIN(M-1,N-1) DO (IF S[M-1] < I THEN (RR:ENDCONS(I,S), : L:CONS(RR,L)))),L);

The function

SUMDOMAIN(arg1 arg2, arg3, arg4)

will, when given the arguments

 $arg_1 = q$, $arg_2 = \sum a_j$, $arg_3 = h$

 $\arg_{4} = [a_1, \ldots, a_q];$

return the list of all lists of indicies

[b₁, ..., b_a]

such that the relations in (3.1) hold:

SUMDOMAIN(M,N,H,L):=BLOCK((D,S),D:DOMAIN(M,N-H), FOR S IN D DO (FOR I THRU M-1 DO (IF NOT (S[I] <= L[I] AND L[I] < S[I+1]) THEN D:DELETE(S,D)), IF S[M] > L[M] THEN D:DELETE(S,D)),D);

The function

FUNDCLASSG(arg1, arg2)

will, when given the arguments

$$\arg_1 = r, \arg_2 = A;$$

generate the fundamental class $\Omega(r + 2, ..., A + 1)$ of G(r, A). For this it uses the function FUNDLIST: ODMAING HR. A. NI :- BLOCK (15, RR LL M. L: CL, M. A. R. L IF N -- MR (M+1)/2 THEN (IF EQUAL (M, I) THEN (IF N -- A+1 THEN L) (M)) ELSE FOR 1 THEN (IF N -- A+1 THEN L) (M) FOR 8 IN DOMAIN(M-1, N+1) (D) (JF 5 (M-1) - 1 THEN (RR:EMDCONS(I.S), L: CONS(FR,L) !)) (L)

The fundtion

SUMDOMAIN (SIG, SIG, SIG, SIG,

will, when given the arguments

ars, - q. arse - Z age ares - b

alga eps] = para

return the list of all lists of indicies

log ever ends

such that the relations in (3.7) hold:

SUMDOMAIN(H, N, H, L):+8LOCK(10.5), D: DOMAIN(H, N-H), FOR S IN D OD (FOR I THRU H-1 00 (JF NOI ISII) -- LII) AND LIII < 8(1+1) THEN D: DELETE (S.D)]. (F S HD > LIM) THEN D: DELETE (S.D)).

> The function FUNDOLASSG(arg₁, arg₂) will, when given the arguments $arg_1 = r$, $arg_2 = A r$ generate the fundamental class O(r + 2, r)

For this it uses the function Full I'

FUNDLIST(R,A):=BLOCK([L],L:[], FOR I FROM R+2 THRU A+1 DO L:ENDCONS(I,L),L);

FUNDCLASSG(R, A): = APPLY(OMEGA, FUNDLIST(R, A));

The function

TOTALDOMAIN(arg1, arg2)

will, when given the arguments

 $\arg_1 = r$, $\arg_2 = A$;

return the list of all possible lists of indicies

$$[a_1, \ldots, a_{A-r}]$$

where

 $1 \le a_1 < a_2 < \ldots < a_{A-r} \le A + 1$:

TOTALDOMAIN(R,A):=BLOCK([S,LO,HI],LO:(A-R)*(A-R+1)/2, HI:SUM(I,I,R+2,A+1),S:[], FOR J FROM LO THRU HI DO S:APPEND(DOMAING(R,A,J),S), S);

The function

```
OMEGATRANSFORM(arg<sub>1</sub>, arg<sub>2</sub>, arg<sub>3</sub>)
```

when given the arguments:

 $\arg_1 = r, \arg_2 = B, \arg_3 = F(c_1(Q), \ldots, c_{r+1}(Q))$, the last one being a polynomial in the Chern-classes of the universal quotient bundle of G(r, A), will apply the substitution (3.1) repeatedly to $F = F \cdot \Omega(r + 2, \ldots, A + 1)$ untill the expression is free of $c_1(Q), \ldots, c_{r+1}(Q)$.

```
FUNDLIST (R. A) - BLOCK ( LL , L: ().
FOR I FROM R.2 THEN ALL DO LIENDONS () . LI.LI:
```

FUNDCLASSE (R. A).1 - APPLY (OHEGA, FUNDLIST (R, A) 1)

The function

COTALDOMAIN(BIS, BIS, BIS,

will, when given the arguments

arg, = r, arg, = A J

return the list of all possible lists of indicies

Lyng and ypage

where

1 1 + 4 4 7-8 > ... > s > 1 = 1

```
TOTAL DOMAINIR. A) 1-BLOCK (IS.LO.HI), LO: (A-RIX (A-R+1)/2,
HIISUN(I.I.R+2.A+1), 3: II.
FOR J FROM LO THRU HI DO 5: APPEND (DOMAING (R.A.J.S), 51,
51:
```

The function

OMBGATRANSFORM (arg, arg, arg,)

when given the arguments:

and $a_{1} = r$, and $a_{2} = 0$, and $a_{3} = 1(c_{1}(0), \cdots, c_{n+1}(0))$, the last one being a polynomial in the Chern-classes of the universal quotient bundle of G(r, A), will apply the substitution (3.1) repeatedly to $F = F \cdot \Omega(r + F, \ldots, A + 1)$ untill the expression is free of $c_{1}(0)$, ..., $c_{n+1}(0)$. OMEGATRANSFORM (R, A, F):=BLOCK ([M, TOT, N, U, OLD, NEW], F:F*FUNDCLASSG (R, A), M:A-R, TOT: TOTALDOMAIN (R, A), FOR I THRU R+1 DO (FOR J THRU INF UNLESS FREEOF (CONCAT (C, I, Q), F) DO. FOR S IN TOT DO (N: SUM (S [K], K, 1, M), U: SUMDOMAIN (A-R, N, I, S), OLD: APPLY (OMEGA, S)*CONCAT (C, I, Q), NEW: SUM (APPLY (OMEGA, U [K]), K, 1, LENGTH (U)), F: RATSUBST (NEW, OLD, F)), F);

It is best to assume that the (weighed) degree of F with respect to the (weighed) degrees of $c_1(Q)$, ..., $c_{r+1}(Q)$ is $\leq \dim = (r + 1)(A - r)$. However, with the usual intersectiontheoretic interpretation of monomials in the c(Q)'s, the function returns the correct result in this case as well, namely zero. On the other hand, the term

 $\deg^2 = (c_1(Q)^{\dim})^2$

which might occur in F (and indeed it will later on) is clearly not intended to mean the element $c_1(Q)$ of A(G(r, A)) raised to the 2dim power, but rather the square of the Chern number $c_1(Q)^{\dim}$. Thus the function OMEGATRANSFORM has to be used with caution in such cases.

The function will for instance return the following results:

(C 11) OMEGATRANSFORM(2, 5, c2Q**2); (D 10) OMEGA(2, 3, 6) + OMEGA(1, 4, 6) or (C 14) OMEGATRANSFORM(2, 5, c1Q**2*c2Q - c2Q**2 - c1Q*c3Q); (D 14) OMEGA(2, 4, 5). The function OMEGATRANSFORM will also convert any polynomial

expression in $c_1(Q)$, ..., $c_{r+1}(Q)$ into the corresponding one involving the Ω 's. Thus for instance we obtain the following: CONECATRANSFORMIA, A.F.) I. BELGEK (EN. TOT.M.U. OLO. NEW) F.F.B.EUNCOLASSG (R. A), R.H.-R. TOT. TOTAL DOTALNIR, AI, FOR I THEO R.I DO FREEGE (CONEAT IC. I. OI, F.I. OO FREEGE (CONEAT IC. I. OI, F.I. OO GOT ST IN TOI DO U.SUTTOONAIMIA-R.M. I.SI, NEW, SUTTOONAIMIA-R.M. I.SI, NEW, SUTTAONAIMIA-R.M. I.SI, NEW, SUTTAFFLY (OTECA. SI & CONCATIC, I. OI F. RATSINGT (NEW, G.O. FIJ), FJ,

It is bast to assume that the (weighed) degree of P with respect to the (weighed) degrees of $c_1(Q)$, ..., $c_{rel}(Q)$ is a dim = (r + 1)(A - r). However, with the usual intersectiontheoretic interpretation of monomials in the c(Q)'s the function returns the correct result in this case as well, manaly zero. Ch the other hand, the term

deg² = (c,(Q)^{dia})²

which might occur in F (and indeed it will later on) is clearly not intended to mean the element $c_1(q)$ of A(G(r, A)) raised to the 2dim power, but rather the square of the Chern mumber $c_1(q)^{dim}$. Thus the function OMEGATHANTION has to be used with coution in such cases.

The function will for instance return the following results:

- C 11) OMEGATRANSFORM(2, 5, c2Q**2) :
- (D 10) OMEGA(2, 3, 6) + OMEGA(1, 4, 6)
 - 10
- (0 14) ONEGATRARSFORM(2, 5, c10**2*c2Q c2Q**2 010*c3Q)) (D 14) ONEGA(2, 4, 5).

The function OMEGATRANSFORM will also convert any polynomial expression in $c_1(Q)$, ..., $c_{r+1}(Q)$ into the corresponding one involving the Ω 's. Thus for instance we obtain the following:

(C 5) OMEGATRANSFORM(1, 3, CHEPO13);

(D 5)
$$6 OMEGA(1,2)T^4 + 12 OMEGA(1,3)T^3 + (7 OMEGA(2,3) + 7 OMEGA(1,4))T^2 + 4 OMEGA(2,4)T + OMEGA(3,4)$$

Here the Chern polynomial of G(1, 3) has previously been computed and assigned to the variable CHEPO13 , using the function G .

Similarly,

(3.

(C 6) OMEGATRANSFORM(1, 3, SEP013);

(D 7)
$$280MEGA(1,2)T^4 - 280MEGA(1,3)T^3 + (90MEGA(2,3) + 90MEGA(1,4))T^2 - 40MEGA(2,4)T + 0MEGA(3,4)$$
.

The results of the function calls with r = 1, A = 4, 5 are given in the appendix, § 3.

What we really need are the degrees of the Chern- and Segre classes, however. To compute these, we may use the following classical formula, given in [HP]:

dim
$$\left(\Omega_{\alpha_{0}}, \ldots, \alpha_{r}\right) = \sum_{i=0}^{r} \alpha_{i} - \frac{1}{2}r(r+1)$$

deg $\left(\Omega_{\alpha_{0}}, \ldots, \alpha_{r}\right) = \left(\frac{\dim \Omega_{\alpha_{0}}, \ldots, \alpha_{r}}{r}\right) = \frac{\left(\dim \Omega_{\alpha_{0}}, \ldots, \alpha_{r}\right)!}{\prod\limits_{i=0}^{r} (\alpha_{i}!)} \prod_{\lambda > \mu} (\alpha_{\lambda} - \alpha_{\mu})$

To get the formulae in a reasonable form, we had to introduce the Schubert-symbols $\Omega_{\alpha_0}, \ldots, \alpha_r$, which are related to $\Omega(a_1, \ldots, a_{A-r})$

(C S) OMEGATRANSFORM(1, 3, CHE2013) ;

(D 5) $6 \text{OMEGA}(1,2)T^2 + 12 \text{OMEGA}(1,3)T^2 + (7 \text{OMEGA}(2,3) + 7 \text{OMEGA}(1,4))T^2 + 4 \text{OMEGA}(2,4)T + 0 \text{OMEGA}(3,4)$

Here the Chern polynomial of G(1, 3) has previously been computed and assigned to the variable CHEPO(3 ; using the function G

(C 6) OMEGATHANSFORM(1, 3, SEPOI3) :

(D.7) $280MEGA(1,2)T^{4} = 280MEGA(1,3)T^{2} + (90MEGA(2,3) + 90MEGA(1,4))T^{4}$ - 40MEGA(2,4)T + 0MEGA(3,4).

The results of the function calls with r = 1, A = 4, 5 are iven in the spendix, § 3.

What we really need are the degrees of the Chern- and Segre classes, however. To compute these, we may use the following classical formula, given in [HP]:

$$a_{4,\alpha} \left(a_{\alpha_0}, \dots, a_{n} \right) = \sum_{k=0}^{n} a_k - b_k (e+1)$$

$$a_{6,\alpha} \left(a_{4,\alpha_0}, \dots, a_{n} \right) = \left(\frac{b_{4,\alpha_0}}{a_{4,\alpha_0}} - \frac{b_{n}(e+1)}{a_{2,\alpha_0}} \right) = \left(a_{1,\alpha_0} - b_{1,\alpha_0} + b_{1,\alpha_0} \right)$$

To get the formulae in a reasonable form, we had to introduce the Schubert-symbols Ω_{α_0} , ..., α_{α_1} , which are related to $\Omega(a_1, \ldots, a_{n-1})$

in the following way (for proof, see [Hm 7] Lemma 3.6):

Lemma 3.3
$$\Omega_{\alpha_0}, \ldots, \alpha_r = \Omega(a_1, \ldots, a_{A-r})$$
, where

$$1 \le a_1 < a_2 < \ldots < a_{A-r} \le A+1$$

are the numbers obtained from

 $\{1, 2, \ldots, A + 1\}$

by deleting

$$\{A - \alpha_{r} + 1, \dots, A - \alpha_{0} + 1\}$$

Thus for instance

$$\Omega_{0,1}, \ldots, r = \Omega(1, \ldots, A - r)$$

and in fact this is the only Schubert cycle of dimension zero. Obviously it has degree 1.

We therefore have an alternative way of computing the degrees of the Chern- and Segre classes: Writing

$$P(t) = 1 + c_1(G)t + ... + c_{dim}(G)t^{dim}$$

where

$$G = G(r, A), dim = (r + 1)(A - r)$$

and c1, ..., cdim are the Chern classes, we put

$$Q(t) = c_1(Q)^{\dim} + c_1(G)c_1(Q)^{\dim-1}t + \cdots$$

+
$$c_{dim-1}(G)c_1(Q)t^{dim-1} + c_{dim}(G)t^{dim}$$

Then putting

$$\omega = \Omega(1, \ldots, A - r) ,$$

OMEGATRANSFORM applied to Q(t) will return

I = a, < a₀, < ... < a_{4-p} = A + 1

are the numbers obtained from

(1, 2, wink + 1)

by deleting

 $(A - a_2 + 1) = A - A_3 + 1)$

Thus for instance

(2 - A x (1) P = 1.... A - 2)

and in fact this is the only Schubert cycle of dimension zero. Obviously it has degree 1.

We therefore have an alternative way of computing the degrees of the Ghern- and Segre classes: Writing

 $P(t) = 1 + e_1(0)t + \dots + e_{qin}(G)t^{dim}$

where

G = G(T, A), dim = (T + 1)(A - T)

and c1. ... catm are the Chern classes, we put

... + 3^{1-mib}(p), p(p), p + ^{mib}(p), p = (3)p

+ cdim-1(0)c, (Q)cdim-1 + cdim(0)cdim

Then putting

 $\omega = \Omega(1 - A \dots A \dots A) \Omega = \omega$

OMEGATRAMSFORM applied to Q(t) will return

 $e_0 \omega + e_1 \omega t + \dots + e_{dim} \omega t^{dim}$

where

$$e_i = deg(c_i(G))$$

Similarly for the Segre Polynomial.

The removal of ω in the output is a matter of stream-lining which I did not carry out. For simplicity I gave the output as "P(t) = Q(t)"

which is of course incorrect. Q(t) is the "degree" of P(t).

Since OMEGATRANSFORM is rather time-consuming, an attempt was made to save some time by arranging the computation somewhat differently, through the function PIERI and POLYSGOMEGA. Their definitions, as well as the results for G(1, 3) and G(1, 4), is given in the appendix, § 3.

Unfortunately the computation became too long already for G(1, 5), so that this approach is clearly not feasible on MACSYMA. However, the principle itself might still work on a computer for higher Grassmanians. What we have demonstrated here, then, is that Schubert Calculus can indeed be implemented on a computer; at least the amount of Schubert Calculus needed for the current project.

Some alternatives to this method will be discussed in Sections 5 and 6.

where

Similarly for the Segre Polynomial.

The removal of ω in the output is a matter of stream-lining which I did not darry out. For simplicity I gave the output as "P(t) = Q(t)"

which is of course incorrect. Q(t) is the "degree" of P(t). Since OMEGATHANSFORM is rather time-consuming. an attempt was made to save some time by arranging the computation schewhat differently, through the function FIRM and POLYSCOMEGA. Their definitions, as well as the results for G(1, 3) and G(1, 4), is given in the annendix. (3.

Unfortunately the computation became too long already for 0(1, 5), so that this approach is clearly not feasible on MACEYMA. However, the principle itself might still work on a computer for higher Grassmanians. What we have demonstrated here, than, is that ichubert calculus can indeed be implemented on a computer; at least the amount of Schubert Calculus needed for the current project.

Some alternatives to this method will be discussed in Sections 5 and 5.

§ 4. Projective embeddings and duality.

Let X be a non-singular, projective variety over the algebraically closed field k, embedded in projective N-space by the embedding

i : X
$$\ensuremath{ \longrightarrow } \mathbb{P}^N_k$$
 .

In this setting, we define the m th embedding-obstruction of the embedded variety X as

$$\gamma_{m} = \deg(X)^{2} - \sum_{j=0}^{m-\dim(X)} {\binom{m+1}{m-\dim(X)-j}} p_{j}(X, i)$$

provided that $m \le 2\dim(X)$, while $\gamma_m = 0$ for $m \ge 2\dim(X) + 1$. Here deg(X) is the degree of X with respect to the embedding i, $\dim(X)$ is the dimension of X and

 $p_j(X, i) = degree(s_j(X))$

is the degree, with respect to the embedding i , of the j th Segre class of X . One then has the following

Theorem 4.1. X may be embedded into \mathbb{P}_k^m by a projection from \mathbb{P}_k^N if and only if $\gamma_m = 0$.

For proofs of this and related results, see for instance [Hm 2, 3, 4, 7], [HR], [Jn], [Lk 2-4] and [Rb 1-5] as well as [K1].

It is sometimes convenient to express γ_m in a slightly different form. In fact, letting D denote the divisor class which corresponds to the embedding i, we have

$$deg(X) = D^{dim(X)}$$

$$p_{j}(X, i) = D^{\dim(X)-j}s_{j}(X)$$
,

where we identify $A^{\dim(X)}(X)$ and Z via the degree-map

§ 4. Projective embeddings and duality.

Let X be a non-singular, projective variety over the algebraically closed field k . embedded in projective H-space by the embedding

In this setting, we define the m th subsiding obstruction of the embedded variety X as

$$Y_{m} = deg(x)^{2} - \sum_{j=0}^{m-dim(x)} {m + l - dim(x) - j} y_{j}(x, x)$$

provided that as 2din(X), while $\gamma_m = 0$ for m > 2din(X) + 1. Here deg(X) is the degree of X with respect to the subcding 1, dim(X) is the dimension of X and

is the degree, with respect to the embedding 1, of the 3 th Segro

Theorem 4.1. X may be embedded into $\mathbf{F}_{\mathbf{K}}^{\mathbf{N}}$ by a projection from $\mathbf{F}_{\mathbf{K}}^{\mathbf{N}}$ if and only if $\mathbf{Y}_{\mathbf{M}} = 0$.

For proofs of this and related results, see for instance [Hm 2, 3, 4, 7], [HM], [Jn], [Lk 2-4] and [Rb 1-5] as well as [X1];

It is sometimes convenient to express yn in a slightly different form. In fact, letting D denote the divisor class which corresponds to the embadding 1, we have

$$p_{3}(X, 1) = D^{dim(X)-j_{a_{3}}(X)}$$

geneergeb and siv & the (X) A viltoebi ew eren

deg:
$$A^{\dim(X)}(X) \longrightarrow Z$$

Thus

$$\gamma_{\mathbf{m}} = \left(p^{\dim(\mathbf{X})} \right)^2 - \sum_{\mathbf{j}=0}^{\mathbf{m}-\dim(\mathbf{X})} \left(\begin{array}{c} \mathbf{m}+1\\ \mathbf{m}-\dim(\mathbf{X}) - \mathbf{j} \end{array} \right) p^{\dim(\mathbf{X})-\mathbf{j}} \mathbf{s}_{\mathbf{j}}(\mathbf{X}) .$$

Furthermore, the generic projection from ${\rm I\!P}^N_k$ to ${\rm I\!P}^m_k$ induces a morphism

$$p_{m}: X \longrightarrow X' \subset \mathbb{P}_{k}^{m},$$

which has a ramification cycle on X denoted by $Ram(p_m)$. We now have (see [Rb], [Jn], [HR]) the

Theorem 4.2. The degree of the cycle $Ram(p_m)$ is given by the expression

$$\operatorname{ram}_{m} = \sum_{j=0}^{m+1-\dim(X)} {\binom{m+1}{j+1}} D^{\dim(X)-j} s_{j}(X)$$

The observation that for $m \leq 2dim(X)$,

$$\gamma_m - \gamma_{m-1} = \operatorname{ram}_{m-1}$$

yields the

<u>Corollary 4.2.1</u>. Assume that $\gamma_{2\dim(X)} = 0$. Then X can be embedded into \mathbb{P}^{m} via a projection from \mathbb{P}^{N} if and only if

$$ram_{m} = 0$$

K. Johnson in [Jn] conjectured that

Thus

$$= \left(c_{1m(x)} \right)^{2} - \sum_{j=0}^{m} c_{1m(x)} \left(\frac{m+1}{m-c_{1m}(x)} - \frac{1}{2} \right)^{2} c_{1m}(x) - \frac{1}{2} c_{2}(x) .$$

Furthermore, the generic projection from R to R Induces a

$$p_{\mathbf{n}}: \quad \mathbf{X} \longrightarrow \mathbf{X}^* \subset \mathbf{X}_{\mathbf{X}}^{\mathbf{n}} \ ,$$

which has a ramification cycle on I denoted by $Ram(p_n)$. We now have (see [Rb], [Jn], [HR]) the

Theorem 4.2. The degree of the cycle Ram(pm) is given by the expression

$$ram_{m} = \sum_{j=0}^{m+1} -\frac{dim(x)}{(j+1)} pdim(x) - i_{s,j}(x)$$

The observation that for a 5 2dim(X) .

yields the

Corollary 4.2.1. Assume that "2dim(X) = 0. Then X can be embedded into Pⁿ via a projection from P^N if and only if

K. Johnson in [Jn] conjectured that

 $\operatorname{ram}_{2\dim(X)-1} = 0 \Rightarrow \gamma_{2\dim(X)} = 0$

This is shown by W. Fulton and J. Hansen in F-H as a corollary of their remarkable <u>connectedness-theorem</u>. One of the initial motivations for the computations below was to check this conjecture against various families of examples, to be generated as described in section 6.

The following functions will generate the Chern polynomial and the Segre polynomial of X in terms of the Chern classes c_1, \ldots, c_{dim} of X:

CHERNPOLY(DIM, T):=1+SUM(CONCAT(C, J)*T^J, J, 1, DIM);

SEGREPOLY (DIM, T) := TAYLOR (1/CHERNPOLY (DIM, T), T, 0, DIM);

We proceed with the functions which return the m th embedding obstruction and the degree of the ramification cycle $Ram(p_m)$:

GAMMASEGRE (DIM, M) := DEG^2 - (BINOMIAL (M+1, M-DIM) *D^DIM +SUM (BINOMIAL (M+1, M-DIM-J) *D^ (DIM-J) *CONCAT (S, J), J, 1, M-DIM)):

It is convenient not to substitute D^{dim} for deg in the first term, see for instance the remark on the use of OMEGATRANSFORM in section 3. Next, we write the same entity in terms of the Chern classes of X:

ram2dim(X)-1 = 0 = 72dim(X) = 0 .

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The following functions will generate the Chern polynomial and the Segre polynomial of X in terms of the Charn classes Cy. Calm of X:

CHERNPOLY (01M, T) = 1 (SUM (CONCATIC, J) # 7*1, J, 1, 01M) 5

SECREPOLY (0111, T) .- TAVLOR () / CHERWOLY (0111, T), T. 0, 0111;

We proceed with the functions which return the n th embedding obstruction and the degree of the ramification cycle Ram (p_n) :

```
GARMASEGRE (010, 1) - DEG^22
- (010001AL (0+1, 0+010) +0*010
+SUM (010001AL (0+1, 0-010-1) +0* (010-1) +CONCATIS, 0) ; J.
1. 0-0101 -1
```

It is convenient not to substitute D' for deg in the first tarm, sea for instance the remark on the use of ONFOATRANSFORM in section 3. Next, we write the same entity in terms of the Chern classes of X:

GAMMACHERN(DIM, M):=BLOCK([Q,P],Q:GAMMASEGRE(DIM,M), P:SEGREPOLY(DIM,T), FOR I THRU DIM DO Q:SUBST(COEFF(P,T,I),CONCAT(S,I),Q),EXPAND(Q));

Analogously we express the degree of the ramification cycle as

follows:

- '

RAMSEGRE (DIM, M) := BINOMIAL (M+1, DIM) *D^DIM +SUM (BINOMIAL (M+1, J+DIM) *D^ (DIM-J) *CONCAT (S, J), J, 1, M-DIM+1);

RAMCHERN(DIM,M):=BLOCK([Q,P],Q:RAMSEGRE(DIM,M), P:SEGREPOLY(DIM,T), FOR I THRU DIM DO Q:SUBST(COEFF(P,T,I),CONCAT(S,I),Q), EXPAND(Q));

Finally we define the m th defect-obstruction of the embedded variety $X \hookrightarrow \mathbb{P}^{\mathbb{N}}$ as

$$\delta_{s} = \sum_{i=s}^{\dim} (i+1)_{s+1} e_{\dim -i} = \sum_{i=s}^{\dim} (i+1)_{s+1} D^{i} c_{\dim -i}(X)$$

The reason for this name is as follows: The dual variety X^{\vee} of X with respect to the embedding i is "normally" a hypersurface, i.e. a subvariety of $\mathbb{P}^{\mathbb{N}}$ of dimension $\mathbb{N} - 1$, see for instance [K1]. But more precisely we have the following result, which is proved in [Hm 7]:

Theorem 4.3. Assume that

$$\delta_0 = \delta_1 = \dots = \delta_{m-1} = 0 , \delta_m \neq 0 .$$

Then

 $\dim(X^{\vee}) = N - 1 - m .$

Following A. Landman [Ln 1, 2] we shall call $m = N - 1 - dim(X^{\vee})$

the (duality) defect of X .

The corresponding function in MACSYMA is:

Tuen

ollowing A. Landman [in 1, 2] we shall call

the (duality) defect of 3

The corresponding function in MACSYMA is:

DELTACHERN(N, DIM, M): = SUM(BINOMIAL(I + 1, M + 1)
 *D^I*CONCAT(C, DIM - I), I, M, DIM - 1)
 + BINOMIAL(DIM + 1, M + 1)*D^DIM

Our intended use of the above functions is the following:

We will consider a certain class of embedded varieties, for which there is given a procedure for computing the Chern classes c_1, \ldots, c_{dim} as well as the intersection numbers

$$D^{n-(i_1+\ldots+i_n)} \stackrel{i_1}{c_1} \cdots \stackrel{i_n}{c_n}, n = \dim$$

This could for instance be the varieties generated from a given one by a finite number of general hypersurface sections, or the ones obtained by blowing up certain loci; possibly both. Reasonable choices for the given variety - which we might call the <u>generating</u> <u>variety</u> - would be a Veronese-variety, or products of some simple given varieties embedded by the Segre-embedding, or finally a Grassmanian or more generally a flag-manifold. We return to this in Section 5.

In an environment established in this way, by a generating variety and a generating procedure, one would then compute the numbers

 $\gamma_{\rm m}$, ram_m, $\delta_{\rm m}$.

We return to the significance of this in Section 5.

The most straightforward approach, which would also be the most flexible and therefore preferable to other alternatives, would be to compute the Chern polynomial of a generated variety, then the intersection numbers above and substitute them into the results of DELTACHERN(N, DIM, N): = SUM(BINOMIAL(I + 1, M + 1)) *D'I* CONCAT(C, DIM - I), I, M, DIM - 1) + BINOMIAL(DIM + 1, M + 1)*D'DIM

Our intended use of the above functions is the following: We will consider a certain class of embedded variaties, for which there is given a procedure for computing the Chern classes

This could for instance be the variables generated from a fiven one by a finite number of general hypersurface sections, or the ches obtained by blowing up certain loci; possibly both. Heasonable choices for the given variaty - which we might call the <u>generating</u> variaty - would be a Veromese-variaty, or products of some simple given variaties embedded by the Segre-embedding, or finally a drassmanian or more generally a flag-manifold. We return to this in Section 5.

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The most straightforward approach, which would also be the most flexible and therefore preferable to other alternatives, would be to compute the Charn polynomial of a generated variaty, then the intersection numbers above and substitute them into the results of RAMSEGRE(DIM, M), RAMCHERN(DIM, M) DELTACHERN(N, DIM, M).

In order to explore the feasability of this, and also to establish the base for using Grassmanians as generating varieties, we carried out some of the above computations for Grassmanians up to G(2, 6). Of course both embedding and duality properties of G(1, A)is well known from geometric theory, see [Hm 7], so no new results were expected from the computations in these cases. In fact, it is known that G(1, A), which is of dimension

$$n = 2(A - 1)$$
,

can be embedded into \mathbb{P}^{2n-3} via a projection from the Plücker embedding, and this is the best possible. Moreover, if A is even and ≥ 4 , then

$$\dim(X^{\vee}) = N - 3$$

where

$$\mathbf{N} = \begin{pmatrix} \mathbf{A} + \mathbf{1} \\ \mathbf{2} \end{pmatrix} - \mathbf{1}$$

is the embedding-dimension of the Plücker embedding. If A is odd, then

 $\dim(X^{\vee}) = N - 1 .$

These results are due to A. Landman, [Lm 1, 2].

However, from what follows we do make the observation that both G(2, 5) and G(2, 6) can not be projected into a space of lower dimension than 2n + 1, where n is the dimension of G(2, 5) resp. G(2, 6), starting with the Plücker embedding. But this can probably be shown easily by direct geometric and elementary methods, and is

RAMSEGRE(DIM. M), RAMOHERN(DIM. M) DELTACHERN(W, DIM. M)

In order to explore the feasability of this, and show to eatablish the base for using Grassmantans as generating variaties, we carried out some of the above computations for Grassmantans up to G(2, 6). Of course both embedding and duality properties of G(1, 4)is well known from geometric theory, see [Hm 7], so no new results were expected from the computations in these cases. In fact, it is known that G(1, 4), which is of dimension

n = 2(A - 1),

can be embedded into P^{-D-2} via a projection from the Fideker embedding, and this is the best possible. Moreover, if A is even

$$\mathcal{E} = \mathbf{X} = (\mathbf{X})_{min}$$

is the embedding-dimension of the Flücker ambedding. If A is odd

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However, from what follows we do make the observation that both G(2, 5) and G(2, 6) can not be projected into a space of lower dimension then 2n + 1, where n is the dimension of G(2, 5) cesp. G(2, 6), starting with the Pificker embedding. But this can probably be shown easily by direct geometric and elementary methods, and is The main point was to see if the computational method referred to above is feasible. Unfortunately the outcome is that at the present time, at least, this direct approach does not seem to be practicable. In sections 5 and 6 we establish an alternative procedure which is less flexible, but which is better suited for large computations.

The function

GRASS(arg₁, arg₂)

will, when given the arguments

 $\arg_1 = r$, $\arg_2 = A$;

print the Chern polynomial of G(r, A) in terms of the Chern classes of Q, and finally evaluate

 $\gamma_{2\dim(X)}$, $\operatorname{ram}_{2\dim(X)-1}$, \cdots

stopping with the first non-zero one. I also worked with a variation of GRASS which always computes at least one ram_i, in order to check Johnson's conjecture. The functions are as follows:

GRASS(R, A):=BLOCK(IDIM, M, GS, DE, N, TEST), DIM: (R+1)*(A-R), CHEPO: CHERNPOLYGRASS(R, A, T), PRINT("THE CHERNPOLYNOMIAL OF GRASS(", R, ", ", A, ") IS: ", CHEPO), PRINT("THE RELATIONS OF THE CHERNCLASSES OF Q ARE:"), RELATIONSOFCHERNCLASSESG(R, A), M: 2*DIM, GS: GAMMASEGRE(DIM, M), PRINT(ARRAYAPPLY(GAMMA, [M]), "=", GS), GC: GAMMACHERN(DIM, M), GC: GRASSEVAL(GC), GC: OMEGATRANSFORM(R, A, GC-DEG^2), DE: PROD(I!, I, 1, R)*DIM!/PROD(I!, I, A-R, A).

```
N:NUMFACTOR(GC)+DE^2,PRINT("= ",N),TEST:0,
FOR I THRU DIM WHILE EQUAL(TEST,0) DO
(M:M-1,GS:RAMSEGRE(DIM,M),
PRINT(ARRAYAPPLY(RAM,[M]),"=",GS),GC:RAMCHERN(DIM,M),
GC:GRASSEVAL(GC),GC:OMEGATRANSFORM(R,A,GC),
N:NUMFACTOR(GC),PRINT("=",N),TEST:N),
PRINT("DEG =",DE));
```

not the main point of what follows.

The main point was to see if and compositional atoms for the to above is feasible. Unfortunately the outcome is that at the present time, at least, this direct approach does not seem to be practicable. In sections 5 and 6 we establish an alternative procedure which is better suited for large

Ing function

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Yedtm(X) Talla (X)-1'

stopping with the first non-zero one. I also worked with a variation of GRABS which always computes at least one ram, , in order to check Johnson's contecture. The functions are as follows:

> CHEPO: CHERNPOL YEARS (R.A.T) CHEPO: CHERNPOL YEARS (R.A.T) PRINT("THE CHERNPOL YEARS (R.A.T) CHEPO: CHEPO: CHEPO: CHEPO: CHEPO: CHERNPOL YEARS (R.A.T) PRINT("THE RELATIONS OF THE CHERNCLASSES OF O ARE."] RELATIONSOF CHERNELASSES (R.A.) RELATIONSOF CHERNELASSES (R.A.) RELATIONSOF CHERNELASSES (R.A.) RELATIONSOF CHERNELASSES (R.A.) CG: CAMMACHERN (DIM.ML, DEC CRASSE VALUE) CG: CHECATRANSFORM (R.A. CE. CECS) CG: CHECATRANSFORM (R.C. CESS) CG: CHECATRANSFORM (R.A. CE. CECS) CG: CHECATRANSFORM (R.A. CE. CECS) CG: CHECATRANSFORM (R.A. CE. CECS) CG: CHECATRANSFORM (R.C. CESS) CG: CHECATRANSFORM (R.C. CESS) CG: CHECATRANSFORM (R.C. CESS) CG: CHECATRANSFORM (R.C. (R.C. (R.C. (R.A. (R.C. N))) CG: CHECATRANSFORM (R.C. (R.C. (R.C. (R.))) CG: CHECATRANSFORM (R.C. (R.C. (R.C. (R.))) CG: CHECATRANSFORM (R. (R. (R.))) CG: CHECATRANSFORM (R. (R. (R.))) CG: CHECATRANSFORM (R. (R. (R.))) CG: CHECATRANSFORM

> > NUTRACTOR (GC), PRINT (""", N), TESTI

The function GRASSEVAL substitutes the Chern-classes for the Grassmanians into the general expressions:

GRASSEVAL (P): =BLOCK ([H], FOR I THRU DIM DO H:RATSUBST (COEFF (CHEPO, T, I), CONCAT (C, I), P), H:RATSUBST (C10, D, H), EXPAND (H), FOR I THRU DIM DO H:RATSUBST (COEFF (CHEPO, T, I), CONCAT (C, I), H), EXPAND (H));

Output from this function is listed in the appendix, § 4. The assertions made above about G(2, 5) and G(2, 6) follow from that.

If one were only interested in the data actually printed out by the function, this would not be the most efficient way to proceed: Indeed, one could then find the degrees of the Segre classes as in section 3, and substitute the result directly into the expressions for $\gamma_{\rm m}$ and $\operatorname{ram}_{\rm m}$. But the function GRASS was also used to give $\gamma_{\rm m}$ and $\operatorname{ram}_{\rm m}$ expressed in terms of $c_1(Q), \ldots, c_{r+1}(Q)$, by inserting a PRINT statement after the statement

GC : GRASSEVAL(FC)

This was done in order to look for a possible pattern, which conceivably could lead to simpler formulae for γ_m and ram_m in terms of the Chern classes of Q. But as no pattern seemed to emerge, it was abandoned. Another approach, which also turned out to be impracticable on MACSYMA, was to determine whether or not γ_m was zero by checking if the expression returned for γ_m in terms of $c_1(Q), \ldots, c_{r+1}(Q)$ was contained in the ideal in $\mathbb{Z}[c_1, \ldots, c_{r+1}]$ generated by the relations of the Chern classes. For instance, we obtained the following results from the last version of GRASS: The function ORASSEVAL substitutes the Chern-classes for the trassmentane into the general expressions:

> ABSEVAL (P) 1-BLOCK (1H). FOR 1 THRU 01H 00 HIRATSUBST (COEFF (CHEPO, T, 1) .CONCAT (C, 11, P). HIRATSUBST (C10, 0, HI. EXPAND (H). FOR 1 THRU DIM DO HIRATSUBST (COEFF (CHEPO, T, 1) .CONCAT (C, 1). H). EXPAND (H) :

Output from this function is listed in the appendix, 9^{+} . The essertions made above about G(2, 5) and G(2, 6) follow from that. If one were only interested in the data actually printed out by the function, this would not be the most efficient way to proceed: Indeed, one could then find the degrees of the Segre classes as in section 5, and substitute the result directly into the expressions for γ_m and range. But the function GASS was also used to give γ_m and range expressed in terms of $c_1(0), \ldots, c_{n+1}(0)$, by inserting a FRIM statement after the statement

This was done in order to look for a possible pattern, which conceivably could lead to simpler formulae for Y₂ and real in terms of the Gnern classes of Q. But as no pattern seemed to smarge, it was abandoned. Another approach, which also turned out to be impracticable on MACSIMA, was to detarmine whether or not Y₂ was zero by checking if the expression returned for Y₂ in terms of c;(Q), ...; c_{r+1}(Q) was contained in the ideal in 3[c₁, ..., c_{r+1}] generated by the relations of the Ghern classes. For instance, we

```
(C9) #rass(1,3);
THE CHERNPOLYNOMIAL OF GRASS( 1 , 3 ) IS:
\begin{pmatrix} 2 & 2 & 4 & 4 & 3 & 3 & 2 & 2 \\ (4 & C2Q & - & 4 & C1Q & C2Q & + & 3 & C1Q & T & + & 6 & C1Q & T & + & 7 & C1Q & T & + & 4 & C1Q & T & + & 1 \\ \end{pmatrix}
 THE RELATIONS OF THE CHERNCLASSES OF Q ARE:
 C1Q - 2 C1Q C2Q = 0
 C2Q - 3 C1Q C2Q + C1Q = 0
                                                  3
 GAMMA = - S4 - 9 D S3 - 36 D S2 - 84 D S1 + DEG - 126 D
 2 2 2 2
= DEG + 4 C2Q - 4 C1Q C2Q - 2 C1Q
= 0
 RAM = S4 + 8 D S3 + 28 D S2 + 56 D S1 + 70 D
 = 4 C1Q C2Q - 4 C2Q
 = 0
 2 3 4
RAM = D S3 + 7 D S2 + 21 D S1 + 35 D
 = 0
 = 0
  \frac{2}{RAM} = D S2 + 6 D S1 + 15 D
  = 0
  = 0
  \frac{3}{RAM} = D S1 + 5 D
  = C1Q
  = 2
  DEG = 2
```

- 36 -

This certainly looks promising: Indeed, both ram_5 and ram_6 vanish identically here. However, it should be pointed out that had the function SEGREPOLYGRASS been used as in GG, then the resulting Chern polynomial would have been more complicated even in this simple case, so that only ram_5 would have vanished identically. We then would have gotten

This certainly looks promising: Indeed, both range and ready vanish identically here. However, it should be pointed out that had the function SEGREPOLYCRASS been used as in GO, then the resulting Chern polynomial would have been more complicated even in this simple case, so that only range would have vanished identically. We then would have gotten

$$RAM_6 = 4C1Q^2C2Q - 2C1Q^4$$

Moreover, we get

$$\gamma_7 = 4c_1(Q)^2 c_2(Q) - 4c_2(Q)^2$$

and since

$$c_1(c_1^3 - 2c_1c_2) - (c_2^2 - c_1^2c_2 - 2c_1^2c_2 + c_1^4)$$

= $c_1^2c_2 - c_2^2$

we also get $\gamma_7 = 0$ by the indicated method.

For G(1, 4) the expressions are more complicated, and in fact none of them vanish identically. But the expressions are still managable enough, so that the method indicated would work. But for G(1, 5), G(2, 5), G(1, 6) and G(2, 6) the expressions become very large, and so do the relations of the Chern classes. Moreover, already the computation of the Chern polynomial beyond G(2, 6) with

the method used is a large affair, and even if such computations could be carried out with disk use and time, it is not clear that the result would be of much use in the given environment. RANK = 4019 029 - 2019

2(2) - 4 - (2) - 5 (2) - 4 - (2)²

and since

or (c12 - sotos) - (c2 - c1es - sotes + c1)

we also get $\gamma_{\gamma} = 0$ by the indicated method:

For G(1, 4) the expressions are more complicated, and in fact none of them vanish identically. But the expressions are still managable anough, so that the method indicated would work. But for G(1, 5), G(2, 5), G(1, 6) and G(2, 6) the expressions become very large, and so do the relations of the Chern classes. Moreover, already the computation of the Chern polynomial beyond G(2, 6) with

the method used is a large affair, and even if such computations could be carried out with disk use and time, it is not clear that the result would be of much use in the given environment.

§ 5. Grassmanians of lines and combinatorical identities.

In this section we show how the numerical study of section 4 may be continued along somewhat different lines from those suggested by the material of section 3, at least in certain special cases.

Indeed, we now take up the case G = G(1, N), Grassmanians of lines in \mathbb{P}^N .

We use the expression

C

$$t(Q \otimes Q^{\vee}) = (4c_2Q - c_1Q^2)t^2 + 1$$

which we computed in section 2 in order to find the Chern and Segre classes of G , by means of (2.4) and (2.5). Actually, using the available information on $c_t(Q \otimes Q^V)$, we may use this method up to Grassmanians of 3-spaces in some \mathbb{P}^N . Of course the expressions then become quite unmanagable, at least "by hand".

Moreover, since Q is of rank 2, and $c_1(Q)$ is the pull back of a hyperplane class via the Plücker embedding, in order to find the degrees of the Chern and Segre classes of G, all we need are the degrees of $c_2(Q)^j$ for $j \leq N - 1$. By Proposition 3.6 of [Hm 7] we now have

$$c_2(Q)^J = \Omega(1, 2, ..., j, j + 3, ..., N + 1)$$
,

and since

$$\{1, 2, ..., j, j + 3, ..., N + 1\} =$$

 $\{1, 2, ..., N + 1\} - \{N - a_1 + 1, N - a_2 + 1\}$

gives

$$N - a_1 + 1 = j + 1$$
, $N - a_0 + 1 = j + 2$

so that

5. Grassmanians of lines and compinatorical identifies.

In this section we show now the numerical study of section 4 may be continued along somewhat different lines from those suggested by

Indeed, we now take up the case G = G(1, H), Grassmanians of lines in \mathbb{R}^{H} .

We use the expression

1 + " + ("P, - P, 0") - (" + 0" P), -

which we computed in section 2 in order to find the Chern and Segre classes of G, by means of (2.4) and (2.5). Actually, using the available information on $c_{\epsilon}(Q Q^{V})$, we may use this method up to Grassmanians of 3-spaces in some P^{R} . Of course the expressions then become quite unmanagable, at least "by hand". Moreover, since Q is of rank 2, and $c_{\epsilon}(Q)$ is the pull back of a

degrees of the Chern and Segre classes of Q, all we need are the degrees of $c_2(Q)^3$ for $j \in \mathbb{N} - 1$. By Proposition 3.5 of [Hm 7] we now have

and since

savia

$$S + 0 = 1 + os - N/(1 + 0 = 1 + ss - 0)$$

so that

- 8e .

$$a_1 = N - j, a_0 = N - j - 1$$

we have that

$$c_2(Q)^j = \Omega_{N-j-1,N-j}$$

by Lemma 3.3. Thus (3.2) gives

(5.1)
$$\deg(c_2(Q)^j) = \frac{(2(N-1-j))!}{(N-1-j)!(N-j)!}$$

where the degree is taken with respect to the Plücker-embedding.

This method may also be used to compute the degrees of the monomials in the Chern classes of Q for higher Grassmanians, i.e. for G(r, N)'s with r > 1. But of course the size of the computations rapidly become quite large, and simple closed form expressions like (5.1) can not be expected.

If the intention with these computations were merely to compute the embedding- and duality numbers of section 4 for Grassmanians, it might not be worth the effort: Indeed, the embedding dimension of G(1, N) via a projection from the Plücker embedding is of course well known and the result easy to prove, [Hm 7] section 3. Further, <u>A. Landman</u> proved in [Lm 1, 2] that the dual variety of G(1, N) with respect to the Plücker embedding is of codimension 3 provided N is even and ≥ 4 , and of codimension 1 otherwise.

However, the main point with these computations is that it becomes possible to use G(1, N)'s - and with the extention indicated above G(r, N)'s with $r \leq 3$ - as generating varieties in the sense of section 6.

Moreover, it turns out as we shall see below, that for G(1, N)'s the above known information on embeddings and duality yields certain

we have that

by Lemma 3.3. Thus (3.2) gives

$$\frac{!((t - t - \pi)s)}{!(t - t - \pi)!} = (L(p)_{s})_{s} = (t, s)$$

where the degree is taken with respect to the Plucker-embedding.

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However, the main point with these computations is that it becomes possible to use G(1, N)'s - and with the extention indicated above G(r, N)'s with $r \le 3$ - as <u>generating variaties</u> in the sense of section 6.

Moreover, it turns out as we shall see below, that for G(1, 3) a the above known information on embeddings and duality yields certain

combinatorical identities, which appear to have no easy direct proofs. Thus the geometry is in this case more likely to yield combinatorical information than the other way around. This is not a new situation in Schubert Calculus.

We start with γ_m and ram_m for G(1, N), and get the following (see [Hm 7] section 4; $c_i = c_i(Q)$):

$$s_{\rho}(G(1,N)) =$$

$$\sum_{j=0}^{\lfloor \frac{\ell}{2} \rfloor} (-1)^{\ell} - j \left\{ \begin{pmatrix} N+\ell-j \\ \ell-j \end{pmatrix} \begin{pmatrix} \ell-j \\ j \end{pmatrix} - 4 \begin{pmatrix} N+\ell-1-j \\ \ell-1-j \end{pmatrix} \begin{pmatrix} \ell-1-j \\ j-1 \end{pmatrix} - \begin{pmatrix} N-\ell-2-j \\ \ell-2-j \end{pmatrix} \begin{pmatrix} \ell-2-j \\ j \end{pmatrix} \right\} c_1^{\ell} - 2^{j} c_2^{j}$$

Hence

$$deg(s_{\ell}(G(1,N))) =$$

$$\begin{bmatrix} \frac{\ell}{2} \end{bmatrix}$$

$$\sum_{\mathbf{j}=0}^{\lfloor \ell-j \rfloor} \begin{pmatrix} \mathbf{N}+\ell-j \\ \ell-j \end{pmatrix} \begin{pmatrix} \ell-j \\ j \end{pmatrix} - 4 \begin{pmatrix} \mathbf{N}+\ell-1-j \\ \ell-1-j \end{pmatrix} \begin{pmatrix} \ell-1-j \\ j-1 \end{pmatrix} - \begin{pmatrix} \mathbf{N}-\ell-2-j \\ \ell-2-j \end{pmatrix} \begin{pmatrix} \ell-2-j \\ j \end{pmatrix} \frac{(2(\mathbf{N}-1-j))!}{(\mathbf{N}-1-j)!(\mathbf{N}-j)!}$$

Letting $n = \dim(G(1, N)) = 2N - 2$, the fact that G(1, N) may be embedded into \mathbb{P}^{2n-3} via a projection from the Plücker embedding is equivalent to

$$\gamma_{2n-3}=0,$$

which after straightforward computations yields the combinatorical identity (for $N \ge 3$):

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We start with γ_{n} and ram_n for Q(1, N), and get the following (see [Hn 7] section 4: $c_{1} = c_{1}(Q)$):

 $\sum_{i=1}^{n} (-1)^{i-1} \left(\binom{n+2-1}{2-2} \binom{n+2-1}{2-2} + \binom{n+1-2}{2-2} \binom{n-2-1}{2-2} \binom{n+2-2-2}{2-2} \binom{n-2-2}{2-2} \binom{n-2-2}{2-2} + \binom{n-2-2}{2-2} \binom{n-2-2}{2$

deg(s,(G(1,N)) =

Letting $n = \dim(G(1, N)) = 2N - 2$, the fact that G(1, N) may be embedded into \mathbb{R}^{2n-3} via a projection from the Flücker embedding is equivalent to

which after straightforward computations yields the combinatorical identity (for M = 3):

$$\sum_{j=0}^{N-3} \sum_{\ell=2j}^{2N-5} (-1)^{\ell-j} \binom{4N-6}{2N-5-\ell} \left\{ \binom{N+\ell-j}{\ell-j} \binom{\ell-j}{j} \right\}$$

$$-4\left(\begin{array}{c}N+\ell-1-j\\\ell-1-j\end{array}\right)\left(\begin{array}{c}\ell-1-j\\j-1\end{array}\right)-\left(\begin{array}{c}N+\ell-2-j\\\ell-2-j\end{array}\right)\left(\begin{array}{c}\ell-2-j\\j\end{array}\right)\frac{\left(2\left(N-1-j\right)\right)!}{\left(N-1-j\right)!\left(N-j\right)!}$$

$$= \left\{ \frac{(2(N-1))!}{(N-1)!N!} \right\}^{2}$$

Similarly one obtains combinatorical identities from $ram_{2n-1} = ram_{2n-2} = ram_{2n-3} = 0$.

We next turn to the duality-deficiency. First we obtain the following expressions for the Chern classes:

 $C_{r}(G(1,N)) =$

$$\sum_{\ell+2s=r}^{\left[\frac{2}{2}\right]} (-1)^{s} (4c_{2}-c_{1}^{2})^{s} \sum_{j=0}^{\left[\frac{N+1}{\ell-j}\right] \binom{\ell-j}{j} c_{1}^{\ell-2j} c_{2}^{j} =$$

0

0≦l≦r

$$\begin{bmatrix} \frac{r}{2} \end{bmatrix} \begin{bmatrix} \frac{r}{2} \end{bmatrix} - s_{s}$$

$$\sum_{s=0} \sum_{j=0}^{r} \sum_{i=0}^{r-2s-j} (-1)^{s+i} + \frac{s-i}{i} + \frac{s-$$

Thus

r rarra

$$deg(e_r(G(1,N))) =$$

$$\sum_{s=0}^{\lfloor z \rfloor \lfloor z \rfloor - s} \sum_{i=0}^{s} (-1)^{s+i} 4^{s-i} {s \choose i} {N+1 \choose r-2s-j} {(r-2s-j) \choose j} \frac{(2(N-1-s-j+i))!}{(N-1-s-j+i)!(N-s-j+i)!}$$

We get

..

$$\delta_{m} = \sum_{i=m}^{n} \binom{i+1}{m+1} deg(c_{n-i}(\Omega_{G(1,N)}^{1})) =$$

$$\sum_{i=m}^{n} \sum_{j=0}^{\lfloor n-i \rfloor} \sum_{k=0}^{\lfloor n-i \rfloor} \sum_{\ell=0}^{j} \sum_{\ell=0}^{j} (-1)^{n-i+j+\ell} \ell_{\mu} j - \ell \binom{i+1}{m+1} \binom{j}{\ell} \binom{N+1}{n-i-j-k} \binom{n-i-j-k}{k-j} \frac{(2(N-1-k+\ell))!}{(N-1-k+\ell)!(N-k+\ell)!}$$

after a straightforward computation.

Computation of the first few values of δ_m yields the table on the following page.

As we see, this is in good agreement with Landmans results quoted above. Moreover, there are many clearly appearing patterns in the table. It would be nice to have <u>geometric</u> proofs for these. So far, we can only ecplain the zeroes, via Landman's results qouted above.

deg(e_r(0(1,N)) =

Jes en

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0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-0-	MZ
	4
- 4 6 6 4 V - 6 W V 4 8 8 6 V 4	ர
2 6 6 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	6
	7
0 0 0 0 0 0 0 0 0 0 0 0 0 0	8
5 20 20 20 20 20 20 20 20 20 20	9
0 0 55 2244 30041 8998704 30244 3932646 678535 377234 4862 4862	10
6 30 150 3-50 3-50 20534 562256 589352 15146352 16036 15146352 16796 1679	11
0 91 91 5824 125502 456144 1421888 3810560 8791016 17456088 2977860 587860 587860 587860 587860	12
7 42 1512 9072 232316 940716 3287396 940716 3287396 2052432 205289829392 306773054 395068198 318938088 205443557178 318938088 205443504 106119712 2288132 2288132 208012	13
0 140 140 12936 79100 1824540 7061920 23825712 70291140 142985825712 2179566460 2179563630 227774029858495 217956363630 3229484170 1967883960 1967883960 196868500 52062432 742900	14
	15

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§ 6. <u>Generating varieties and classes of</u> projective varieties.

A well known conjecture by Horrocks and Mumford, [HM], asserts that if X is a variety embedded in \mathbb{P}^n , and if X is non-singular and of codimension 2, then X is a complete intersection provided $n \ge 6$.

A related conjecture by Hartshorne, [Hn] asserts the same conclusion under the assumption that X is nonsingular and of dimension $> \frac{2}{3}$ n, provided $n \ge 7$.

These conjectures clearly testify to the current lack of information on examples of non singular projective varieties of high dimension.

If a classification theory of the type carried out by Swinnerton-Dyer [Sw] for varieties of degree 4 could be carried out in general, these and many other questions would of course be settled. But such a general classification-theory is clearly not in sight at this time.

However, the fundamental idea underlying Swinnerton-Dyers classification is to obtain all varieties from a fundamental set of varieties by processes such as blowing up subvarieties and taking hyperplane sections. While such a procedure might not yield a complete classification in general, it still can be used to generate lists of examples, and provide the starting point for a systematical search for counterexamples to conjectures such as the above.

Initially one should search for subvarieties of projective space, of high diemnsion and of relatively low codimension. One would then be able to get some idea of how frequent non-singular varieties of this type are, and of course examine such questions as whether they are complete intersections or not.

The procedure which we propose is therefore to start out with a variety such as for instance a suitable Grassmanian, a Veronese-

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The procedure which we propose is therefore to start out with a variety such as for instance a suitable Grassmanian, a Veronesevariety or a product of projective spaces embedded by the Segreembedding (see [Hm 7], [Da] and [Vi] for details on the last two cases). We would then modify the situation by blowing up suitable subschemes, such as a finite number of possibly mutiple points or Schubert-varieties in the case of Grassmanians. In all of these cases there are well known algorithms for computation of the Chern- and Segre polynomials of the new variety in terms of that of the old one and data associated with the center of blowing up. Moreover, the new variety comes with a natural projective embedding, given by the projective embedding of the blow-up of the ambient space with the given center induced from the interpretation of a blow-up as a monoidal transformation.

Thus there are natural projective embeddings of the results of each modification, and furthermore algorithms for the computation of the corresponding degrees of the Chern- and Segre classes.

Another possible modification is for instance to take sections with hypersurfaces, in which case the degrees of the new Chern- and Segre classes are expressed in terms of the corresponding data for the original variety and the degree of the hypersurface by particularly simple formulae.

Thus once an initial variety is selected - which we will call the <u>generating variety</u> - a whole class of projective varieties is generated by all possible modifications of the type described above applied repeatedly in any order. Of course the class of <u>admissible</u> <u>modifications</u> may vary, one could for example start with the simplest one which is to take hypersurface sections. Also, it is no essential limitation to restrict the permissible zero dimensional centers for the blowing ups to simple points.

To get a good picture of each such generated class of projective varieties, one should ideally have classified the varieties in them according to projective equivalence. In principle this should be

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- 44

possible to do on a computer, with the availability of a sufficiently powerfull system for symbolic manipulation like MACSYMA. But given the results from section 3 in particular, I rather doubt that an investigation along these lines is feasible at this time. This is mainly due to the size such computations would necessarily have.

However, there is a coarser equivalence of a homological nature which is well suited for a computer. Under this equivalence varieties for which the degrees of all monomials in the Chern classes are equal, and for which the Chow rings are not "too different", will be identified.

All of the above adresses itself to non-singular varieties. But singular varieties may also be considered. In the singular case we will have to compute the invariants introduced in [Hm 4], and which are actually degrees of different types of Segre classes introduced later in the singular case by Fulton and Mc Phersson. Moreover, one also needs the degrees of Fultons singular [FM]. Chern classes, [F] and [Hm 6]. If one allows singularities, the computation of the invariants of the modified variety in terms of those of the original one becomes more difficult, and further development of the general theory is needed before this can be done. The reason why this would be of interest, is among other things the following: We may obtain more non singular varieties if we allow singular generating varieties. Also, the choice of admissible modifications is greatly expanded, since we may take cones over given subvarieties, or more generally form the join of two given projective subvarieties of some projective space, see [AK]. Furthermore, we can modify by deformations as in [Hm 5] or as in [Hm 1].

The program outlined above will constitute the continuation of this paper, [Hm 8] and [Hm 9].

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- 46 -Appendix Printouts

§ 1. Output from the functions TENSOR and EXTERIOR 2.

```
(C12) A[1];
                                  C1AT + 1
([]12)
(C13) A[2];
                                   2
                              C2AT + C1AT + 1
([13)
(C14) A[3];
                                       2
                               3
                          C3AT + C2AT + C1AT + 1
(D14)
(C15) B[1];
                                   C1BT+1
(D15)
(C16) B[2];
                               C2B T + C1B T + 1
(D16)
(C17) B(3];
                               3
                                        2
                          C3B T + C2B T + C1B T + 1
(D17)
(C18) TENSOR (A[1],1,B[1],1);
                               (C1B + C1A) T + 1
(D18)
(C19) TENSOR (A[1],1,B[2],2);
                                    2
                                        2
                (C2B + C1A C1B + C1A ) T + (C1B + 2 C1A) T + 1
(D19)
(C20) TENSOR (A[1],1,B[3],3);
                                                                          2
                                      3
                                         3
(D20) (C3B + C1A C2B + C1A C1B + C1A ) T + (C2B + 2 C1A C1B + 3 C1A ) T
                                                          + (C1B + 3 C1A) T + 1
 (C21) TENSOR (A [2], 2, B [2], 2);
 (D21) (C2B - 2 C2A C2B + C1A C1B C2B + C1A C2B + C2A + C1B C2A
 + C1A C1B C2A) T + (2 C1B C2B + 2 C1A C2B + 2 C1B C2A + 2 C1A C2A + C1A C1B
                                                          2
                                                              2
                                        2
       2
                 + (2 C2B + 2 C2A + C1B + 3 C1A C1B + C1A ) T
  + C1A C1B) T
  + (2 C1B + 2 C1A) T + 1
 (C22) TENSOR (A [3], 3, B [3], 3);
                                        2
                                                         2
 (D22) (C3B + 3 C3A C3B + C1A C2B C3B - 2 C1B C2A C3B - 3 C1A C2A C3B
                                      2.
                        3
                 2
                             2
  + C1A C1B C3B + C1A C3B + 3 C3A C3B - 3 C1B C2B C3A C3B
  - CIA C2B C3A C3B - C1B C2A C3A C3B - 3 C1A C2A C3A C3B - 2 C1A C1B C3A C3B
                                 2
   - 2 CIA CIB C3A C3B + C2A C2B C3B - 2 C2A C2B C3B + CIA C1B C2A C2B C3B
```

. . Antenut from the functions TENSON and EXTERIOR C.

2 + CIA C2A C2B C3B + C2A C3B + C1B C2A C3B + C1A C1B C2A C3B + C3A - 3 C1B C2B C3A - 2 C1A C2B C3A + C1B C2A C3A + C1B C3A + C1A C1B C3A 3 2 2 2 2 2 + C2B C3A - 2 C2A C2B C3A + C1A C1B C2B C3A + C1A C2B C3A + C2A C2B C3A + C1B C2A C2B C3A + C1A C1B C2A C2B C3A) T 2 .2: + (3 C2B C3B - 6 C2A C3B + 2 C1A C1B C3B + 3 C1A C3B - 3 C2B C3A C3B - 3 C2A C3A C3B - 6 C1B C3A C3B - 14 C1A C1B C3A C3B - 6 C1A C3A C3B + 2 C1A C2B C3B - 3 C1A C2A C2B C3B + 2 C1A C1B C2B C3B + 2 C1A C2B C3B + 2 C1B C2A C3B + 3 C1A C2A C3B + 2 C1A C1B C2A C3B + 2 C1A C1B C2A C3B - 6 C2B C3A + 3 C2A C3A + 3 C1B C3A + 2 C1A C1B C3A + 3 C1B C2B C3A + 2 CIA C2B C3A - 3 CIB C2A C2B C3A + 2 CIA CIB C2B C3A + 2 C1A C1B C2B C3A + 2 C1B C2A C3A + 2 C1B C2A C3A + 2 C1A C1B C2A C3A 2 + C2A C2B - 2 C2A C2B + C1A C1B C2A C2B + C1A C2A C2B + C2A C2B 2 2 2 8 + C1B C2A C2B + C1A C1B C2A C2B) T + (3 C1B C3B + 3 C1A C3B - 21 C1B C3A C3B - 21 C1A C3A C3B + 3 C2B C3B - 6 C2A C2B C3B + 6 C1A C1B C2B C3B + 6 C1A C2B C3B + 3 C2A C3B + 3 C1A C1B C2A C3B + 3 C1A C2A C3B + 2 C1A C1B C3B + 2 C1A C1B C3B + 3 C1B C3A + 3 C1A C3A + 3 C2B C3A - 6 C2A C2B C3A + 3 C1B C2B C3A + 3 CIA CIB C2B C3A + 3 C2A C3A + 6 CIB C2A C3A + 6 CIA CIB C2A C3A + 2 C1A C1B C3A + 2 C1A C1B C3A + C1A C2B + 3 C1B C2A C2B 2 2 3 2 2 2 2 + C1A C1B C2B + C1A C2B + 3 C1A C2A C2B + 3 C1A C1B C2A C2B

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- 84

3 3 2 + 3 CIA CIB C2A C2B + CIB C2A + CIB C2A + CIA CIB C2A) T + (3 C3B - 21 C3A C3B + 6 C1B C2B C3B + 8 C1A C2B C3B - 4 C1B C2A C3B 2 2 + 3 C1A C2A C3B + 4 C1A C1B C3B + 8 C1A C1B C3B + 2 C1A C3B + 3 C3A + 3 C1B C2B C3A - 4 C1A C2B C3A + 8 C1B C2A C3A + 6 C1A C2A C3A + 2 C1B C3A 2 2 3 2 + 8 C1A C1B C3A + 4 C1A C1B C3A + C2B + C2A C2B + 4 C1A C1B C2B 3 C1A C2B + C2A C2B + 4 C1B C2A C2B + 7 C1A C1B C2A C2B 2 3 2 + 4 C1A C2A C2B + 2 C1A C1B C2B + 2 C1A C1B C2B + C2A + 3 C1B C2A + 4 CIA CIB C2A + 2 CIA CIB C2A + 2 CIA C1B C2A) T + (6 C2B C3B - 3 C2A C3B + 3 C1B C3B + 10 C1A C1B C3B + 6 C1A C3B - 3 C2B C3A + 6 C2A C3A + 6 C1B C3A + 10 C1A C1B C3A + 3 C1A C3A + 3 C1B C2B + 5 C1A C2B + 6 C1B C2A C2B + 6 C1A C2A C2B + 5 C1A C1B C2B + 8 C1A C1B C2B + 2 C1A C2B + 5 C1B C2A + 3 C1A C2A + 2 C1B C2A 2 2 2 2 3 3 2 5 + 8 C1A C1B C2A + 5 C1A C1B C2A + C1A C1B + C1A C1B) T + (6 C1B C3B + 6 C1A C3B + 6 C1B C3A + 6 C1A C3A + 3 C2B + 3 C2A C2B + 3 C1B C2B + 12 C1A C1B C2B + 6 C1A C2B + 3 C2A + 6 C1B C2A 2 2 . 3 + 12 C1A C1B C2A + 3 C1A C2A + 2 C1A C1B + 5 C1A C1B + 2 C1A C1B) T + (3 C3B + 3 C3A + 6 C1B C2B + 7 C1A C2B + 7 C1B C2A + 6 C1A C2A + C1B + 7 C1A C1B + 7 C1A C1B + C1A) T + (3 C2B + 3 C2A + 3 C1B + 8 C1A C1B + 3 C1A) T + (3 C1B + 3 C1A) T + 1 (C23) EXTERIOR\2(A[1],1); Part fell off end. (C24) EXTERIOR\2(A[2],2);

. .

(D24) C1A T + 1 (C25) EXTERIOR\2(A[3],3);

(D25) = C3A T + C1A T + 2 C1A T + 1

(C26) EXTERIOR\2(A[4],4);



```
§ 2. Output from the functions G, GG and RE.
(C5) G(3);
+ 4 C10 T + 1
SEPO 1 3 = (-4 C20 + 4 C10 C20 + 14 C10) T - 14 C10 T
                                                   3 3
                                            . 2. 2
                                        + 9 C10 T - 4 C10 T + 1
C10 - 2 C10 C20 = 0
2 2 4
C2Q - 3 C1Q C2Q + C1Q = 0
(C6) G(4);
                     2
                                             66
                              2
CHEPO 1 4 = (4 C20 + 12 C10 C20 - 13 C10 C20 + 4 C10) T
2 3 5 5
+ (20 C10 C20 - 20 C10 C20 + 10 C10<sup>-</sup>). T
                                      5
 2 2 4 4 3 3
+ (4 C2Q - 7 C1Q C2Q + 14 C1Q ) T + 15 C1Q T
                                          3 3
2 2
+ (C2Q + 11 C1Q) T + 5 C1Q T + 1
2 2 2 4. 6 6
SEPO 1 4 = (3 C2Q + 79 C1Q C2Q - 426 C1Q C2Q + 198 C1Q ) T
 2 3 5 5
+ (5 C1Q C2Q + 150 C1Q C2Q - 105 C1Q ) T'
                                   5 5
 . 3 3
 + (14 C1Q - C2Q) T - 5 C1Q T + 1
           2
c_{20}^{2} - 3 c_{10}^{2} c_{20}^{2} + c_{10}^{2} = 0
2 3 5.
. 6
 - C2Q + 6 C1Q C2Q - 5 C1Q C2Q + C1Q = 0
 (C7) G(5);

    4
    2
    3
    4
    2
    6

    CHEPO
    1
    5
    =
    (9
    C20
    -
    18
    C10
    C20
    +
    42
    C10
    C20
    -
    26
    C10
    C20

 8 8 3 2 5 7 7
+ 5 C1Q ) T + (75 C1Q C2Q - 60 C1Q C2Q + 15 C1Q ) T
  3 2 2 4 6 6
+ (- 6 C20 + 78 C10 C20 - 66 C10 C20 + 25 C10 ) T
```

a outnut from the functions 0, 60 and RE .

- 15 -

2 3 5 5 + (30 C10 C20 - 30 C10 C20 + 30 C10) T 4 4 2 2 + (7 C20 - 2 C10 C20 + 31 C10) T + (6 C10 C20 + 26 C10) T + (2 C2Q + 16 C1Q) T + 6 C1Q T + 1 SEPO 1 5 = (- 52 C2Q + 1318 C1Q C2Q - 828 C1Q C2Q - 902 C1Q C2Q 8 8 3 3 2 5 + 500 C10) T + (- 252 C10 C20 + C10 C20 + 904 C10 C20 7 7 3 2 2 4 6 - 455 C1Q) T + (26 C2Q + 18 C1Q C2Q - 522 C1Q C2Q + 319 C1Q) 6 2 3 5 5 T + (6 C10 C20 + 230 C10 C20 - 194 C10) T 3 3 3 2 2 4 4 + (- 3 C2Q - 78 C1Q C2Q + 105 C1Q) T + (18 C1Q C2Q - 50 C1Q) T + (20 C10 - 2 C20) T - 6 C10 T + 1 CHEPO 2 5 = (8 C3Q - 24 C1Q C2Q C3Q + 4 C1Q C3Q + 52 C1Q C2Q C3Q - 36 C10 C20 C30 + 7 C10 C30 - 8 C10 C20 - 20 C10 C20 5 2 7 9 9 + 38 C1Q C2Q - 21 C1Q C2Q + 4 C1Q) T 2 2 2 3 5 + (- 22 C10 C30 + 56 C10 C20 C30 - 12 C10 C20 C30 + 2 C10 C30 3 2 4 C2Q - 48 C1Q C2Q + 85 C1Q C2Q - 57 C1Q C2Q + 14 C1Q) T + (- 30 C10 C30 + 24 C20 C30 + 36 C10 C20 C30 + 2 C10 C30 - 24 C1Q C2Q + 70 C1Q C2Q - 78 C1Q C2Q + 28 C1Q) T. 3 + (- 15 C30 + 30 C10 C20 C30 + 26 C10 C30 + 43 C10 C20 2 4 6 6 2 2 - 84 C10 C20 + 43 C10) T + (41 C10 C30 + 22 C10 C20 2.

- 52 -

- .63 C10 C20 + 50 C10) T + (24 C10 C30 + 3. C20 - 27 C10 C20 + 45 C1Q) T + (6 C3Q - 6 C1Q C2Q + 32 C1Q) T + 17 C1Q T + 6 C10 T + 1 SEPO 2 5 = (- 404 C30 + 1212 C10 C20 C30 - 9364 C10 C30 2 2 - 1750 C10 C20 C30 + 19266 C10 C20 C30 - 9684 C10 C30 4 3 3 5 2 7 - 70 C10 C20 + 1082 C10 C20 - 10860 C10 C20 + 10572 C10 C20 9 9 2 2 2 2 - 2586 C1Q) T + (2596 C1Q C3Q - 8 C1Q C2Q C3Q - 5184 C10 C20 C30 + 3878 C10 C30 + 13 C20 - 18 C10 C20 4 2 6 8 8 + 2923 C10 C20 - 4200 C10 C20 + 1271 C10) T + (- 510 C10 C30 + 12 C20 C30 + 1008 C10 C20 C30 - 1530 C10 C30 3 3 2 5 7 7 - 12 C10 C20 - 558 C10 C20 + 1590 C10 C20 - 632 C10) T + (51 C30 - 102 C10 C20 C30 + 586 C10 C30 + 35 C10 C20 - 570 C10 C20 + 322 C10) T + (- 197 C10 C30 + 14 C10 C20 + 183 C10 C20 - 168 C10) T + (48 C10 C30 - 3 C20 - 45 C10 C20 3 + 88 C1Q) T + (- 6 C3Q + 6 C1Q C2Q - 44 C1Q) T + 19 C1Q T - 6 C1Q T + 1 2 2 C10 C30 + C20 - 3 C10 C20 + C10 = 03 - 2 C2Q C3Q + 3 C1Q C3Q + 3 C1Q C2Q - 4 C1Q C2Q + C1Q = 0 C30 - 6 C10 C20 C30 + 4 C10 C30 - C20 + 6 C10 C20 - 5 C10 C20 + C10= 9 2 **3** C10 C30 + 3 C20 C30 - 12 C10 C20 C30 + 5 C10 C30 - 4 C10 C20 + 10 C10 C20 - 6 C10 C20 + C10 = 0 2 - 3 C20 C30 + 6 C10 C30 + 12 C10 C20 C30 - 20 C10 C20 C30

- 53 -

- 22 -

5 4 2 3 4 2 6 8 + 6 C10 C30 + C20 - 10 C10 C20 + 15 C10 C20 - 7 C10 C20 + C10

= 0 3 2 · 3 2 3 2 2 C3Q - 12 C1Q C2Q C3Q + 10 C1Q C3Q - 4 C2Q C3Q + 30 C1Q C2Q C3Q 4 6 4 3 3 - 30 C10 C20 C30 + 7 C10 C30 + 5 C10 C20 - 20 C10 C20 5 2 7 9 + 21 C1Q C2Q - 8 C1Q C2Q + C1Q = 0 (C8) G(6): 5 2 4 4 3 6 2 CHEPO 1 6 = (9 C2Q + 36 C1Q C2Q - 102 C1Q C2Q + 106 C1Q C2Q - 43 C1Q C2Q + 6 C1Q) T + (63 C1Q C2Q - 168 C1Q C2Q 3 · 3 5 2 7 9 9 + 252 C10 C20 - 126 C10 C20 + 21 C10) T 9: 9 4 2 3 4 2 6 8 8 + (3 C2Q - 87 C1Q C2Q + 310 C1Q C2Q - 194 C1Q C2Q + 41 C1Q) T 3 2 + (- 42 C10 C20 + 259 C10 C20 - 182 C10 C20 + 56 C10) T 3 2 2 4 6 6 + (C2Q + 127 C1Q C2Q - 93 C1Q C2Q + 62 C1Q) T 6 6 2 3 5 5 + (49 C10 C20 - 14 C10 C20 + 63 C10) T 2 2 4 4 3 3 + (9 C2Q + 20 C1Q C2Q + 57 C1Q) T + (14 C1Q C2Q + 42 C1Q) T 2 2 + (3 C2Q + 22 C1Q) T + 7 C1Q T + 1 SEPO 1 6 = (- 1414 C10 C20 - 18914 C10 C20 + 47362 C10 C20 8 10 10 - 33264 C1Q C2Q + 6916 C1Q) T 4 3 3 5 2 7 + (700 C10 C20 + 2254 C10 C20 - 14448 C10 C20 + 14616 C10 C20 - 3808 C1Q) T + (- 78 C2Q + 162 C1Q C2Q + 4218 C1Q C2Q
 6
 8
 3
 3
 2

 - 6618 C10
 C20
 + 2196 C10
 T
 + (- 154 C10 C20
 - 1120 C10
 C20
 2 5 7 7 3 2 2 + 2954 C10 C20 - 1274 C10) T + (26 C20 + 236 C10 C20

-1228 c10⁴ c20 + 716 c10⁶) 1⁶ + (-28 c10 c20² + 448 c10³ c20 5 5 4 2 4 - 378 C1Q) T + (182 C1Q - 133 C1Q C2Q) T + (28 C10 C20 - 77 C10) T + (27 C10 - 3 C20) T - 7 C10 T + 1 2 2 C20 + 6 C10 C20 - 5 C10 C20 + C10 = 0- C2Q + 15 C1Q C2Q - 35 C1Q C2Q + 28 C1Q C2Q - 9 C1Q C2Q + C1Q = 03 3 3 CHEPO 2 6 = (8 C30 + 8 C10 C20 C30 + 48 C10 C30 + 16 C20 C30 2 2 2 4 2 6 2 - 6 C10 C20 C30 - 150 C10 C20 C30 + 47 C10 C30 4 3 3 5, 2 - 6 C1Q C2Q C3Q - 40 C1Q C2Q C3Q + 211 C1Q, C2Q C3Q - 132 C10 C20 C30 + 23 C10 C30 + 9 C20 - 45 C10 C20 4 4 6 3 8 2 10 + 105 C10 C20 - 170 C10 C20 + 125 C10 C20 - 41 C10 C20 12 12 3 2 3 2 2 + 5 C1Q) T + (8 C2Q C3Q + 114 C1Q C3Q + 90 C1Q C2Q C3Q 3 2 5 2 4 2 3 - 414 C10 C20 C30 + 86 C10 C30 - 9 C20 C30 - 102 C10 C20 C30 . + 575 C10 C20 C30 - 345 C10 C20 C30 + 62 C10 C30 - 18 C10 C20 3 4 5 3 7 2 9 + 120 C10 C20 - 387 C10 C20 + 364 C10 C20 - 141 C10 C20 11 11 3 2 2 2 2 2 + 20 C10) T + (56 C10 C30 + 63 C20 C30 - 259 C10 C20 C30 4 2 3 3 2 - 77 C1Q C3Q - 56 C1Q C2Q C3Q + 539 C1Q C2Q C3Q - 259 C10 C20 C30 + 56 C10 C30 - 15 C20 + 123 C10 C20 6 2 - 516 C10 C20 + 564 C10 C20 - 262 C10 C20 + 45 C10) T

- 55

- 55 -

+ (- 7 c_{30}^{-3} - 28 c10 c20 c30⁻² - 266 c10³ c30² + 22 c20³ c30 2 2 4 6 + 230 C10 C20 C30 + 47 C10 C20 C30 + 29 C10 C30 + 51 C10 C20 3 3 5 2 7 9 5 - 375 C10 C20 + 537 C10 C20 - 337 C10 C20 + 75 C10) T + (- 15 C20 C30 - 247 C10 C30 + 66 C10 C20 C30 + 185 C10 C20 C30 5 4 2 3 4 2 + 64 C10 C30 + 22 C20 - 185 C10 C20 + 395 C10 C20 - 362 C10 C20 + 106 C10) T + (- 105 C10 C30 - 15 C20 C30 + 122 C1Q C2Q C3Q + 143 C1Q C3Q - 44 C1Q C2Q + 232 C1Q C2Q - 322 C10 C20 + 127 C10) T + (- 9 C30 + 20 C10 C20 C30 + 169 C10 C30 - 11 C20 + 115 C10 C20 - 221 C10 C20 + 128 C10) 5 5 2 2 4 4 + 109 C10) T + (42 C10 C30 + 6 C20 - 25 C10 C20 + 80 C10) T + $(7 \ C30 + 49 \ C10)$ T + $(C20 + 23 \ C10)$ T + $7 \ C10$ T + 1SEPO 2 6 = (36.99 C30 - 12142 C10 C20 C30 + 3105 C10 C30 3 2 2 2 2 4 22 + 4758 C20 C30 - 68043 C10 C20 C30 + 131775 C10 C20 C30 - 137452 C10 C30 - 8338 C10 C20 C30.+ 197449 C10 C20 C30 - 480172 C10 C20 C30 + 442503 C10 C20 C30 - 102743 C10 C30 6 2 5 4 4 6 3 + 603 C20 - 738 C10 C20 - 121479 C10 C20 + 344794 C10 C20 - 354873 C10 C20 + 142370 C10 C20 - 19272 C10) T + (528 C20 C30 - 11341 C10 C30 + 12100 C10 C20 C30 + 914 C1Q C2Q C3Q + 26340 C1Q C3Q + 77 C2Q C3Q - 40312 C10 C20 C30 + 107674 C10 C20 C30 - 116728 C10 C20 C30

- 68

8 5 3 4 5 3 + 28184 C10 C30 + 882 C10 C20 + 29844 C10 C20 - 103008 C10 C20 7 2 9 11 11 + 117192 C10 C20 - 49076 C10 C20 + 6644 C10) T 9 11 11 3 2 2 2 2 2 + (3696 C10 C30 - 1188 C20 C30 - 6509 C10 C20 C30 4 2 3 3 2 + 281 C10 C30 + 5734 C10 C20 C30 - 21074 C10 C20 C30 5 7 5 2 4 + 20790 C10 C20 C30 - 3252 C10 C30 - 125 C20 - 5617 C10 C20 4 3 6 2 8 10 + 28028 C10 C20 - 33579 C10 C20 + 13005 C10 C20 - 1322 C10) 10 3 2 3 2 3 T + (- 462 C30 + 1358 C10 C20 C30 - 3206 C10 C30 - 376 C20 C30 2 2 4 6 + 4219 C10 C20 C30 + 339 C10 C20 C30 - 2616 C10 C30 + 641 C10 C20 - 7085 C10 C20 + 8032 C10 C20 - 1367 C10 C20 9 9 2 2 2 2 2 - 515 C10) T + (- 66 C20 C30 + 1667 C10 C30 - 824 C10 C20 C30 - 2035 C10 C20 C30 + 2579 C10 C30 - 25 C20 + 1601 C10 C20 - 1538 C10 C20 - 1219 C10 C20 + 885 C10) T + (- 462 C10 C30 + 90 C20 C30 + 808 C10 C20 C30 - 1481 C10 C30 3 3 2 5 7 7 - 268 C10 C20 + 274 C10 C20 + 1134 C10 C20 - 748 C10) T 2 3 3 2 2 + (58 C30 - 146 C10 C20 C30 + 657 C10 C30 + 22 C20 - 84 C10 C20 4 6 6 2 - 620 C10 C20 + 506 C10) T + (8 C20 C30 - 229 C10 C30 2 3 5 5 + 30 C10 C20 + 254 C10 C20 - 298 C10) T + (56 C10 C30 - 5 C20 - 76 C10 C20 + 155 C10) T + (- 7 C30 + 14 C10 C20 - 70 C10) T + (26 C10 - C20) T - 7 C10 T + 1

2 3 3 2 2 4 C30 - 6 C10 C20 C30 + 4 C10 C30 - C20 + 6 C10 C20 - 5 C10 C20 + C1Q = 0 2 2 2 4 3 3 C1Q C3Q + 3 C2Q C3Q - 12 C1Q C2Q C3Q + 5 C1Q C3Q - 4 C1Q C2Q $3^{\circ} 2$ 5 $2^{\circ} + 10^{\circ} C10^{\circ} C20^{\circ} - 6^{\circ} C10^{\circ} C20^{\circ} + C10^{\circ} = 0^{\circ}$ - 3 C2Q C3Q + 6 C1Q C3Q + 12 C1Q C2Q C3Q - 20 C1Q C2Q C3Q
 5
 4
 2
 3
 4
 2
 6
 8

 + 6
 C10
 C30
 +
 C20
 10
 C10
 C20
 +
 15
 C10
 C20
 7
 C10
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 C10

 3
 2
 3
 2
 3
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 2

 C30
 12
 C10
 C20
 C30
 +
 10
 C10
 C30
 4
 C20
 C30
 +
 30
 C10
 C20
 C30
 4 6 4 3 3 - 30 C10 C20 C30 + 7 C10 C30 + 5 C10 C20 - 20 C10 C20 5 2 7 9 + 21 C10 C20 - 8 C10 C20 + C10 = 0 3 2 2 2 2 2 4 2 4 C1Q C3Q + 6 C2Q C3Q - 30 C1Q C2Q C3Q + 15 C1Q C3Q - 20 C10 C20 C30 + 60 C10 C20 C30 - 42 C10 C20 C30 + 8 C10 C30 5 2 4 4 3 6 2 8 - C2Q + 15 C1Q C2Q - 35 C1Q C2Q + 28 C1Q C2Q - 9 C1Q C2Q 10 + C1Q - 4 C2Q C3Q + 10 C1Q C3Q + 30 C1Q C2Q C3Q - 60 C1Q C2Q C3Q 5 2 4 2 3 4 2 + 21 C10 C30 + 5 C20 C30 - 60 C10 C20 C30 + 105 C10 C20 C30
 6
 8
 5
 3
 4

 - 56
 C10
 C20
 C30
 +
 9
 C10
 C30
 6
 C10
 C20
 +
 35
 C10
 C20
 5 3 7 2 9 1 - 56 C10 C20 + 36 C10 C20 - 10 C10 C20 + C10 4 3 3 3 3 2 C30 - 20 C10 C20 C30 + 20 C10 C30 - 10 C20 C30 2 2 2 4 2 6 2 + 90 C10 C20 C30 - 105 C10 C20 C30 + 28 C10 C30 4 3 3 5 2 + 30 C10 C20 C30 - 140 C10 C20 C30 + 168 C10 C20 C30 7 9 6 2.5 4 4 - 72 C10 C20 C30 + 10 C10 C30 + C20 - 21 C10 C20 + 70 C10 C20 - 84 $C1Q^{6}C2Q^{3} + 45 C1Q^{8}C2Q^{2} - 11 C1Q^{10}C2Q + C1Q^{12}$

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- 84 C19 029 + 45 C19 C29 - 11 C19 C29+ C19

[DSK, USERS] (D4)(C5) LINEL:70; 70 (D5) (C6) GG(3); 2 2 4 4 NEWCHEPO 1 3 = (6 C2Q - 16 C1Q C2Q + 8 C1Q) T NEWSEPO 1 3 = (- 6 C2Q - 16 C1Q C2Q + 25 C1Q) T 3 3 2 2 + (4 C1Q C2Q - 16 C1Q) T + 9 C1Q T - 4 C1Q T + 1 (DG) DONE (C7) GG(4):
 3
 2
 2
 4
 6
 6

 NEWCHEPO
 1
 4
 =
 (14
 C2Q
 +
 88
 C1Q
 C2Q
 72
 C1Q
 C2Q
 +
 16
 C1Q
)
 T
 6 6 + (30 C10 C20 - 40 C10 C20 + 16 C10) T 2 2 4 4 3 3 + (6 C20 - 13 C10 C20 + 16 C10) T + 15 C10 T . 2 2 + (C2Q + 11 C1Q) T + 5 C1Q T + 1 3 2 2 2 NEWSEPO 1 4 = (25 C2Q - 15 C1Q C2Q - 245 C1Q C2Q + 140 C1Q) T 2 3 5 5 + (15 C10 C20 + 110 C10 C20 - 91 C10) T 2 2 4 4 3 3 + (- 5 C2Q - 40 C1Q C2Q + 55 C1Q) T + (10 C1Q C2Q - 30 C1Q) T 3 3 + (14 C10 - C20) T - 5 C10 T + 1(D7)DONE (C8) GG(5); 4 2 3
 4
 2
 3
 4
 2
 6

 NEWCHEPO
 1
 5
 =
 (47
 C2Q
 368
 C1Q
 C2Q
 +
 504
 C1Q
 C2Q
 224
 C1Q
 C2Q
 8 3 3 2 + 32 C10) T + (- 84 C10 C20 + 248 C10 C20 - 160 C10 C20 7 7 3 2 2 4 6 6 + 32 C1Q) T + (- 8 C2Q + 105 C1Q C2Q - 96 C1Q C2Q + 32 C1Q) T 2 3 5 5 + (36 C10 C20 - 38 C10 C20 + 32 C10) T 5 5 2 3 3 + (7 C20 - 2 C10 C20 + 31 C10) T + (6 C10 C20 + 26 C10) T

+ $(2 C2Q + 16 C1Q^2) T^2 + 6 C1Q T + 1$ NEWSEPO 1 5 = (- 98 C2Q + 560 C1Q C2Q + 1134 C1Q C2Q
 6
 8
 8
 3
 3
 2

 - 2436
 C10
 C20
 +
 825
 C10
)
 T
 +
 (168
 C10
 C20
 336
 C10
 C20
 5 7 7 3 2 2 + 1260 C10 C20 - 540 C10) T + (28 C20 + 63 C10 C20 4 6 6 3 5 5 - 588 C1Q C2Q + 336 C1Q) T + (238 C1Q C2Q - 196 C1Q) T 2 2 4 4 3 3 + (- 3 C2Q - 78 C1Q C2Q + 105 C1Q) T + (18 C1Q C2Q - 50 C1Q) T + (20 C10 - 2 C20) T - 6 C10 T + 1 NEWCHEPO 2 5 = (- 142 C3Q + 1062 C1Q C2Q C3Q - 1032 C1Q C3Q - 162 C20 C30 + 198 C10 C20 C30 - 36 C10 C20 C30 + 88 C10 C30 4 3 3 5 2 7 + 648 C10 C20 - 2058 C10 C20 + 2090 C10 C20 - 876 C10 C20 2 2 + 128 C1Q) T + (132 C2Q C3Q - 393 C1Q C3Q + 162 C1Q C2Q C3Q
 3
 5
 4
 2
 3

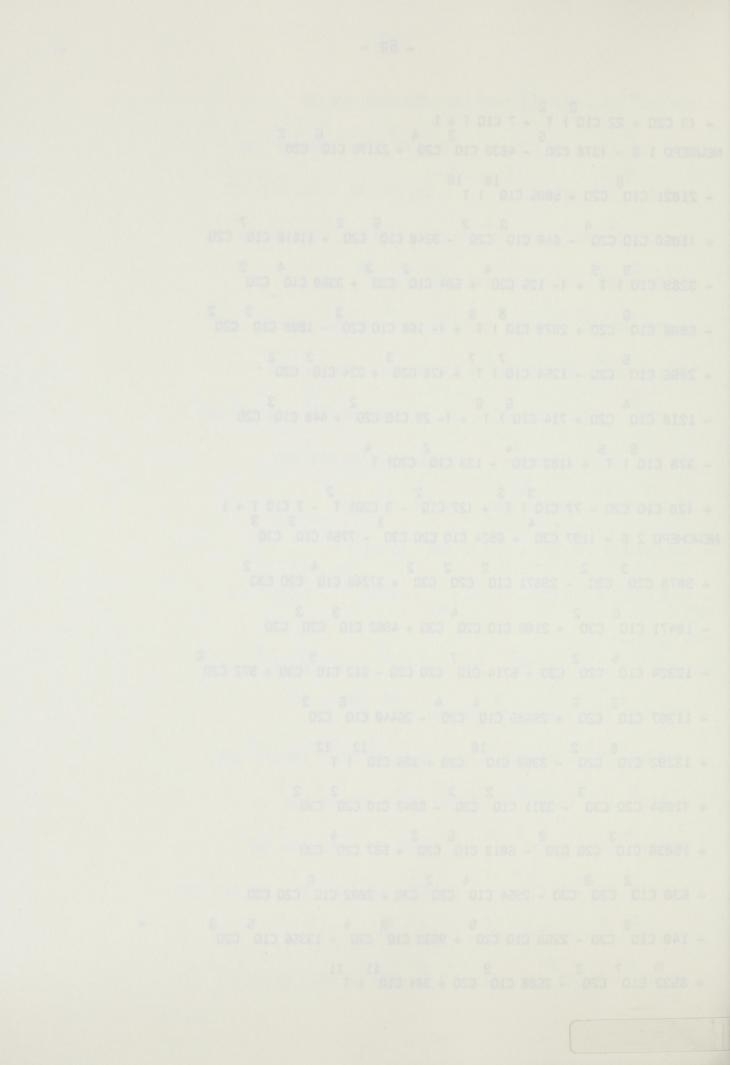
 - 174
 C10
 C20
 C30
 +
 140
 C10
 C30
 +
 81
 C20
 684
 C10
 C20
 4 2 6 8 8 + 1078 C1Q C2Q - 612 C1Q C2Q + 112 C1Q) T + (- 102 C10 C30 + 42 C20 C30 - 120 C10 C20 C30 + 144 C10 C30 3 3 2 5 7 7 - 162 C10 C20 + 474 C10 C20 - 396 C10 C20 + 96 C10) T + (- 12 C30 - 42 C10 C20 C30 + 116 C10 C30 - 20 C20 2 2 4 6 E + 166 C10 C20 - 228 C10 C20 + 80 C10) T + (- 6 C2Q C3Q + 72 C1Q C3Q + 42 C1Q C2Q - 108 C1Q C2Q + 64 C1Q) + (6 C30 - 6 C10 C20 + 32 C10) T + 17 C10 T + 6 C10 T + 1 3 2 . 3 2 NEWSEPO 2 5 = (- 218 C30 + 1746 C10 C20 C30 - 3720 C10 C30

- 60 -

- 150 c20³ c30 - 666 c10² c20² c30 + 5502 c10⁴ c20 c30 - 3416 C1Q C3Q + 420 C1Q C2Q - 1610 C1Q C2Q - 364 C1Q C2Q 7 9 9 2 2 2 + 2016 C1Q C2Q - 670 C1Q) T + (- 204 C2Q C3Q + 1281 C1Q C3Q - 18 C10 C20 C30 - 1998 C10 C20 C30 + 1890 C10 C30 - 45 C20 2 3 6 8 8 + 580 C10 C20 - 1176 C10 C20 + 489 C10) T + (- 330 C1Q C3Q + 30 C2Q C3Q + 564 C1Q C2Q C3Q - 954 C1Q C3Q 3 3 2 5 7 7 - 150 C1Q C2Q + 90 C1Q C2Q + 630 C1Q C2Q - 344 C1Q) T + (48 C30 - 102 C10 C20 C30 + 424 C10 C30 + 20 C20 - 70 C10 C20 6 - 300 C10 C20 + 231 C10) T + (6 C20 C30 - 156 C10 C30 2 3 5 5 + 30 C10 C20 + 120 C10 C20 - 146 C10) T + (42 C1Q C3Q - 6 C2Q - 36 C1Q C2Q + 85 C1Q) T 3 3 2 2 + (- 6 C3Q + 6 C1Q C2Q - 44 C1Q) T + 19 C1Q T - 6 C1Q T + 1 (D8)(C9) GG(6): 5 2 NEWCHEPO 1 6 = (- 135 C2Q + 1532 C1Q C2Q - 2936 C1Q C2Q 6 2 8 10 1 + 2064 C10 C20 - 608 C10 C20 + 64 C10) T 10 10 + (329 C10 C20 - 1176 C10 C20 + 1232 C10 C20 - 480 C10 C20 9 9 4 2 3 4 2 + 64 C1Q) T + (39 C2Q - 347 C1Q C2Q + 656 C1Q C2Q 6 8 8 3 3 2 - 352 C10 C20 + 64 C10) T + (- 56 C10 C20 + 315 C10 C20 2 - 224 C10 C20 + 64 C10) T + (- C20 + 139 C10 C20 - 103 C10 C20 6 6 2 3 5 5 + 64 C1Q) T + (49 C1Q C2Q - 14 C1Q C2Q + 63 C1Q) T + (9 C20 + 20 C10 C20 + 57 C10) T + (14 C10 C20 + 42 C10) T

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2 2 + (3 C2Q + 22 C1Q) T + 7 C1Q T + 1 NEWSEPD 1 6 = (378 C2Q - 4830 C1Q C2Q + 22176 C1Q C2Q 8 10 10 - 21021 C1Q C2Q + 5005 C1Q) T 4 3 3 5 2 7 + (1050 C10 C20 - 840 C10 C20 - 9240 C10 C20 + 11616 C10 C20 9 9 4 2 3 4 2 - 3289 C1Q) T + (- 126 C2Q + 504 C1Q C2Q + 3360 C1Q C2Q 6 8 8 3 3 2 - 6006 C10 C20 + 2079 C10) T + (- 168 C10 C20 - 1008 C10 C20 5 7 7 3 2 2 + 2856 C1Q C2Q - 1254 C1Q) T + (28 C2Q + 224 C1Q C2Q ' 4 6 6 2 3 - 1218 C1Q C2Q + 714 C1Q) T + (- 28 C1Q C2Q + 448 C1Q C2Q 5 5 4 2 4 - 378 C10) T + (182 C10 - 133 C10 C20) T + (28 C10 C20 - 77 C10) T + (27 C10 - 3 C20) T - 7 C10 T + 1 3 3 3 NEWCHEPO 2 6 = (197 C30 + 6924 C10 C20 C30 - 7764 C10 C30 3 2 2 2 2 4 2 + 3078 C20 C30 - 29871 C10 C20 C30 + 37245 C10 C20 C30 6 2 4 3 3 - 10471 C10 C30 + 2106 C10 C20 C30 + 4962 C10 C20 C30 5 2 7 9 6 - 12324 C10 C20 C30 + 5714 C10 C20 C30 - 612 C10 C30 + 972 C20 - 11367 C10 C20 + 26685 C10 C20 - 26440 C10 C20 8 2 10 12 12 + 13292 C1Q C2Q - 3360 C1Q C2Q + 336 C1Q) T + (1064 C20 C30 - 3311 C10 C30 - 6048 C10 C20 C30 3 2 5 2 4 + 15036 C10 C20 C30 - 6013 C10 C30 + 567 C20 C30 - 630 C10 C20 C30 - 2954 C10 C20 C30 + 2002 C10 C20 C30 8 5 3 4 5 3 - 140 C10 C30 - 2268 C10 C20 + 9639 C10 C20 - 13356 C10 C20 7 2 9 11 1 + 8532 C10 C20 - 2608 C10 C20 + 304 C10) T 11 11



3 2 2 2 2 + (- 994 C10 C30 - 582 C20 C30 + 4791 C10 C20 C30 - 3063 C1Q C3Q - 648 C1Q C2Q C3Q - 12 C1Q C2Q C3Q 5 7 5 2 4 + 266 C10 C20 C30 + 156 C10 C30 - 243 C20 + 2781 C10 C20 4 3 6 2 8 10 1 - 5976 C1Q C2Q + 5100 C1Q C2Q - 1952 C1Q C2Q + 272 C1Q) T + (- 154 C30 + 1050 C10 C20 C30 - 1309 C10 C30 - 140 C20 C30 + 406 C10 C20 C30 - 294 C10 C20 C30 + 308 C10 C30 + 567 C10 C20 3 3 5 2 7 9 5 - 2268 C1Q C2Q + 2772 C1Q C2Q - 1392 C1Q C2Q + 240 C1Q) T 2 2 2 2 3 + (114 C20 C30 - 423 C10 C30 + 204 C10 C20 C30 - 286 C10 C20 C30 5 4 2 3 4 2 + 348 C10 C30 + 61 C20 - 680 C10 C20 + 1324 C10 C20
 6
 8
 8
 2
 2

 - 928 C1Q
 C2Q
 + 208 C1Q
 T
 + (- 84 C1Q C3Q
 + 42 C2Q
 C3Q
 2 4 3 2 - 126 C10 C20 C30 + 308 C10 C30 - 140 C10 C20 + 532 C10 C20 - 560 C10 C20 + 176 C10) T + (- 6 C30 - 24 C10 C20 C30 3 3 2 2 4 6 + 220 C10 C30 - 14 C20 + 172 C10 C20 - 288 C10 C20 + 144 C10) 6 2 2 3 5 5 T + (119 C10 C30 + 42 C10 C20 - 112 C10 C20 + 112 C10) T + (42 C10 C30 + 6 C20 - 25 C10 C20 + 80 C10) T 3 3 2 2 + (7 C3Q + 49 C1Q) T + (C2Q + 23 C1Q) T + 7 C1Q T + 1

 4
 3
 3
 3

 NEWSEPO 2 6 = (966 C30
 - 18312 C10 C20 C30
 + 39592 C10
 C30

 - 1484 C20 C30 + 53396 C10 C20 C30 - 156618 C10 C20 C30 + 85932 C10 C30 - 5544 C10 C20 C30 - 6552 C10 C20 C30 + 114828 C10 C20 C30 - 140448 C10 C20 C30 + 41272 C10 C30

- 63 -

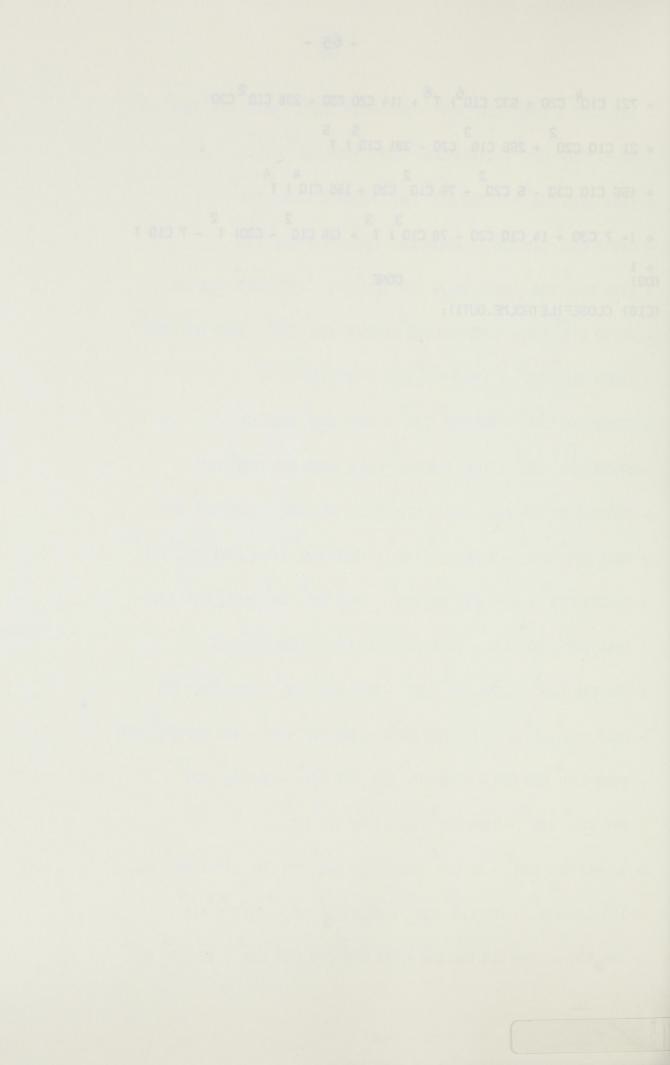
-294 c2a⁶ + 10332 c1a² c2a⁵ - 35700 c1a⁴ c2a⁴ + 14322 c1a⁶ c2a³ 8 2 10 12 1 + 33957 C1Q C2Q - 28028 C1Q C2Q + 5551 C1Q) T 3 2 3 2 2 + (1848 C2Q C3Q - 11256 C1Q C3Q - 8820 C1Q C2Q C3Q 3 2 5 2 4 + 51632 C10 C20 C30 - 40530 C10 C30 + 686 C20 C30 2 3 4 2 6 - 1736 C1Q C2Q C3Q - 38178 C1Q C2Q C3Q + 68712 C1Q C2Q C3Q 8 5 3 4 5 3 - 25179 C1Q C3Q - 1764 C1Q C2Q + 12432 C1Q C2Q - 9240 C1Q C2Q 7 2 9 11 1 - 16830 C1Q C2Q + 17842 C1Q C2Q - 4082 C1Q) T 3 2 2 2 2 2 + (2352 C10 C30 + 756 C20 C30 - 13776 C10 C20 C30 + 17262 C1Q C3Q + 1064 C1Q C2Q C3Q + 10136 C1Q C2Q C3Q 5 7 5 2 4 - 30660 C1Q C2Q C3Q + 14616 C1Q C3Q + 154 C2Q - 3472 C1Q C2Q 4 3 6 2 8 10 1 + 4998 C10 C20 + 7518 C10 C20 - 10857 C10 C20 + 2926 C10) T 3 2 3 2 3 + (- 273 C30 + 2646 C10 C20 C30 - 6412 C10 C30 - 196 C20 C30 2 2 4 6 - 1841 C1Q C2Q C3Q + 12124 C1Q C2Q C3Q - 7980 C1Q C3Q 4 3 3 5 2 7 + 686 C10 C20 - 2240 C10 C20 - 2898 C10 C20 + 6252 C10 C20 - 2035 C1Q) T + (- 273 C2Q C3Q + 1960 C1Q C3Q + 126 C1Q C2Q C3Q - 4060 C1Q C2Q C3Q + 4032 C1Q C3Q - 70 C2Q + 791 C1Q C2Q 4 2 6 8 8 + 882 C10 C20 - 3360 C10 C20 + 1365 C10) T + (- 441 C10 C30 + 21 C20 C30 + 1064 C10 C20 C30 - 1841 C10 C30 - 196 C10 C20 - 161 C10 C20 + 1652 C10 C20 - 876 C10) T 2 3 3 2 2 + (55 C30 - 186 C10 C20 C30 + 732 C10 C30 + 25 C20 - 15 C10 C20

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· (55 C30 - 186 C10 C20 C30 + 732 C10 C30 + 25 C20 - 15 C10 C20

```
-721 \ C10^{4} \ C20 + 532 \ C10^{6}) \ T^{6} + (14 \ C20 \ C30 - 238 \ C10^{2} \ C30
+ 21 \ C10 \ C20^{2} + 266 \ C10^{3} \ C20 - 301 \ C10^{3}) \ T^{5}
+ (56 \ C10 \ C30 - 5 \ C20^{2} - 76 \ C10^{2} \ C20 + 155 \ C10^{3}) \ T^{4}
+ (-7 \ C30 + 14 \ C10 \ C20 - 70 \ C10^{3}) \ T^{3} + (26 \ C10^{2} - C20) \ T^{2} - 7 \ C10 \ T
+ (0) DONE
```

```
(C10) CLOSEFILE (HOLME, OUT1);
```



(D2) [DSK, USERS]
<pre>(C3) GRIND(RE); RE(R):=BLOCK((F,G,H,K,L,M,N),F:X^(R+1)+SUM(CONCAT(C,I,Q)*X^(R+1-I),I,1,R+1), G:RATSUBST(X-1,X,F),H:F-G,K:RESULTANT(F,H,X),L:K, FOR I THRU R+1 DO L:RATSUBST(CONCAT(C,I,Q)*T^I,CONCAT(C,I,Q),L), M:L*COEFF(L,T,0),N:SUM(EXPAND(COEFF(M,T,I))*T^I,I,0,(R+1)^2),N)\$ (D3) DONE</pre>
(C4) RE (1); 2 2
(D4) $(4 C20 - C10) T + 1$
(C5) RE [2]; .
CONCAT FASL DSK MAXOUT being loaded loading done
$\begin{array}{cccccccccccccccccccccccccccccccccccc$
$\begin{array}{cccccccccccccccccccccccccccccccccccc$
(CG) RE [3]; 3 2 2 2 2 2 2
(D6) (256 C40 - 192 C10 C30 C40 - 128 C20 C40 + 144 C10 C20 C40
4 2 2 2 2 2 - 27 C1Q C4Q + 144 C2Q C3Q C4Q - 6 C1Q C3Q C4Q - 80 C1Q C2Q C3Q C4Q
3 + 18 C10 C20 C30 C40 + 16 C20 C40 - 4 C10 C20 C40 - 27 C30
3 3 3 3 2 2 2 2 12 + 18 C1Q C2Q C3Q - 4 C1Q C3Q - 4 C2Q C3Q + C1Q C2Q C3Q) T
+ (- 192 C20 C40 + 72 C10 C40 + 216 C30 C40 - 120 C10 C20 C30 C40
3 3 2 2 3 + 18 C10 C30 C40 + 32 C20 C40 - 6 C10 C20 C40 - 54 C10 C30
2 2 2 2 4 2 3 + 18 C2Q C3Q + 42 C1Q C2Q C3Q - 9 C1Q C3Q - 26 C1Q C2Q C3Q
3 2 5 2 4 10 + 6 C10 C20 C30 + 4 C20 - C10 C20) T
+ (- 112 C40 + 56 C10 C30 C40 + 24 C20 C40 - 32 C10 C20 C40 + 6 C10 C40
+ 48 C20 C30 - 25 C10 C30 - 54 C10 C20 C30 + 38 C10 C20 C30 - 6 C10 C30
+ 17 C2Q - 12 C1Q C2Q + 2 C1Q C2Q) T
+ (16 C2Q C4Q - 6 C1Q C4Q + 26 C3Q - 30 C1Q C2Q C3Q + 8 C1Q C3Q + 28 C2Q

2 2 4 6 E - 24 C1Q C2Q + 8 C1Q C2Q - C1Q) T 6 6 2 4 2 + (8 C4Q - 2 C1Q C3Q + 22 C2Q - 16 C1Q C2Q + 3 C1Q) T 2 2 + (8 C2Q - 3 C1Q) T + 1 (C7) RE[4]; You have run out of LIST space. Do you want more? Type ALL; NONE; a level-no. or the name of a space. ALL: You have run out of LIST space. Do you want more? Type ALL; NONE; a level-no. or the name of a space. ALL: You have run out of LIST space. Do you want more? Type ALL; NONE; a level-no. or the name of a space. ALL; You have run out of FIXNUM space. Do you want more? Type ALL; NONE; a level-no. or the name of a space. ALL: 115235 msec. 139662 msec. 170254 msec. 202918 msec. max allocation exceeded FIXNUM storage capacity exceeded 203907 msec. so far (C8) CLOSEFILE (HOLME, OUT2);

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2029977 2000 20 520

. (C8) CLOSEFILE (HOLHE, OUT2) 1

§ 3. Output from Schubert Calculus. (C5) GRIND(PIERI); PIERI (R, A, TOT, F) := BLOCK ([M, N, U, OLD, NEW], F:F*FUNDCLASSG(R,A),M:A-R, FOR I THRU R+1 DO (FOR J THRU INF UNLESS FREEOF (CONCAT (C, I, Q), F) DO FOR S IN TOT DO (N:SUM(S[K],K,1,M),U: SUMDOMAIN (A-R, N, I, S), OLD: APPLY (OMEGA, S) *CONCAT (C, I, Q), NEW: SUM (APPLY (OMEGA, U[K]), K, 1, LENGTH(U)). F:RATSUBST (NEW, OLD, F))), F)\$ (D5) DONE (C6) GRIND (POLYSGOMEGA); POLYSGOMEGA(R, A):=BLOCK([DIM, CH, D, SE], DIM: (R+1)*(A-R), CH: CHERNPOLYG (R, A, T), TOT: TOTALDOMAIN (R, A), PRINT ("CHEPO", R, A, "=", CH), H: CH, H:PIERI(R,A,TOT,H),PRINT("= ",H), H:C1Q^DIM*RATSUBST(T/C1Q, T, CH), H: PIERI (R, A, TOT, H), PRINT ("= ", H), D: 1/CH. SE: TAYLOR (D. T. Ø, DIM), SE:1+SUM(EXPAND(COEFF(SE,T,I))*T^I,I,1,DIM), PRINT ("SEPO", R, A, "=", SE), E: PIERI (R, A, TOT, SE), PRINT("= ",E),E:C1Q^DIM*RATSUBST(T/C1Q,T,SE), E:PIERI (R, A, TOT, E), PRINT ("= ", E), PRINT ("DONE"))\$ (DG) DONE (C7) LINEL:70: 70 (D7) (C8) POLYSGOMEGA(1,3); CONCAT FASL DSK MAXOUT being loaded loading done HAYAT FASL DSK MACSYM being loaded loading done 2 2 3 3 2 2 CHEPO 1 3 = (4 C2Q - 4 C1Q C2Q + 3 C1Q) T + 6 C1Q T + 7 C1Q T+ 4 C10 T + 13 6 OMEGA(1, 2) T + 12 OMEGA(1, 3) T (7 OMEGA(2, 3) + 7 OMEGA(1, 4)) T + 4 OMEGA(2, 4) T + OMEGA(3, 4) 2 6 OMEGA(1, 2) T + 12 OMEGA(1, 2) T + 14 OMEGA(1, 2) T + 8 OMEGA(1, 2) T + 2 OMEGA(1, 2) 2 3 3 SEPO 1 3 = (- 4 C2Q + 4 C1Q C2Q + 14 C1Q) T - 14 C1Q T

+ 3 C10² T²- 4 C10 T +

 $+ 9 C10^{2} T^{2} - 4 C10 T + 1$ 28 OMEGA(1, 2) T - 28 OMEGA(1, 3) T 2 + (9 OMEGA(2, 3) + 9 OMEGA(1, 4)) T - 4 OMEGA(2, 4) T + OMEGA(3, 4) 3 28 OMEGA(1, 2) T - 28 OMEGA(1, 2) T + 18 OMEGA(1, 2) T - 8 OMEGA(1, 2) T + 2 OMEGA(1, 2) DONE DONE (D8) (C9) POLYSGOMEGA(1,4): 2 2 CHEPO 1 4 = (4 C2Q + 12 C1Q C2Q - 13 C1Q C2Q + 4 C1Q) T 5 5 2 + (20 C10 C20 - 20 C10 C20 + 10 C10) T 2 2 4 4 3 3 + (4 C2Q - 7 C1Q C2Q + 14 C1Q) T + 15 C1Q T + (C2Q + 11 C1Q) T + 5 C1Q T + 1 10 OMEGA(1, 2, 3) T + 30 OMEGA(1, 2, 4) T + (35 OMEGA(1, 3, 4) + 25 OMEGA(1, 2, 5)) T + (15 OMEGA(2, 3, 4) + 30 OMEGA(1, 3, 5)) T + (11 OMEGA(2, 3, 5) + 12 OMEGA(1, 4, 5)) T + 5 OMEGA(2, 4, 5) T + OMEGA (3, 4, 5) You have run out of LIST space. Do you want more? Type ALL; NONE; a level-no. or the name of a space. ALL; 6 = 10 OMEGA(1, 2, 3) T + 30 OMEGA(1, 2, 3) T + 60 OMEGA(1, 2, 3) T + 75 OMEGA(1, 2, 3) T + 57 OMEGA(1, 2, 3) T + 25 OMEGA(1, 2, 3) T + 5 OMEGA(1, 2, 3) SEPO 1 4 = (3 C2Q + 79 C1Q C2Q - 426 C1Q C2Q + 198 C1Q) T 3 5 5 + (5 C1Q C2Q + 150 C1Q C2Q - 105 C1Q) T 2 2 3 3 + (- 3 C2Q - 46 C1Q C2Q + 57 C1Q) T + (10 C1Q C2Q - 30 C1Q) T

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+ (14 C10² - C20) T² - 5 C10 T + 1 63503 msec. = 220 OMEGA(1, 2, 3) T⁶ - 220 OMEGA(1, 2, 4) T⁵ + (125 OMEGA(1, 3, 4) + 65 OMEGA(1, 2, 5)) T⁴ + (- 30 OMEGA(2, 3, 4) - 50 OMEGA(1, 3, 5)) T³ + (14 OMEGA(2, 3, 5) + 13 OMEGA(1, 4, 5)) T² - 5 OMEGA(2, 4, 5) T + OMEGA(3, 4, 5) 83168 msec. = 220 OMEGA(1, 2, 3) T⁶ - 220 OMEGA(1, 2, 3) T⁵ + 190 OMEGA(1, 2, 3) T⁶ - 130 OMEGA(1, 2, 3) T³ + 68 OMEGA(1, 2, 3) T² - 25 OMEGA(1, 2, 3) T + 5 OMEGA(1, 2, 3) DONE DONE

(C10) CLOSEFILE (HOLME, OUT4);

```
La CLO<sup>2</sup> - COD T<sup>2</sup> - SCD T<sup>2</sup> - SCD T+1
CORECALL, 2, 3) T<sup>2</sup> - 223 ORECALL, 2, 4) T<sup>2</sup>
LOS ORECALL, 3, 4) + 65 ORECAL, 2, 50 T<sup>2</sup>
LOS ORECAL, 3, 4) - 56 ORECAL, 2, 50 T<sup>2</sup>
Lo ORECAL, 3, 5) - 13 ORECAL, 5, 5) T<sup>2</sup> - 5 ORECAL, 5, 5) T<sup>2</sup>
A DECALS, 4, 5)
A DECALS, 4, 5)
A DECALS, 4, 5)
A DECALS, 2, 5) T - 220 ORECAL, 2, 5) T<sup>2</sup> - 5 ORECALS, 5, 5)
A DECALS, 4, 5)
A DECALS, 4, 5)
A DECALS, 5, 5) T - 220 ORECAL, 5, 5) T - 5 ORECALS, 5, 5)
A DECALS, 4, 5)
A DECALS, 5, 5) T - 220 ORECAL, 2, 5) T - 5 ORECALS, 5, 5)
A DECALS, 5, 5) T - 220 ORECAL, 2, 5) T - 5 ORECALS, 5, 5)
```

```
(CLO) CLOSEFILE (HOLME, OUTA);
```

§ 4. Output from Grass.

```
2 2
                                             2
                                                     2
                                                                4
                                                                            3 3
                                                                   4
THE CHERNPOLYNOMIAL OF GRASS( 1 , 3 ) IS: (4 C20 - 4 C10 C20 + 3 C10 ) T + 6 C10 T + 7 C10 T + 4 C10 T + 1
THE RELATIONS OF THE CHERNCLASSES OF Q ARE:
  3
C10 - 2 C10 C20 = 0
  2
           2
C2Q - 3 C1Q C2Q + C1Q
                        = 8
                                     3
                                              2
                           2
GAMMA = - S4 - 9 D S3 - 36 D S2 - 84 D S1 + DEG - 126 D
    8
= 8
RAM = S4 + 8 D S3 + 28 D S2 + 56 D S1 + 70 D
  7
- 8
                2
                          3
    = D S3 + 7 D S2 + 21 D S1 + 35 D
RAM
   6
. 8
        2
                 3
RAM = D S2 + 6 D S1 + 15 D
  5
= 8
        3
RAM = D S1 + 5 D
  4
= 2
DEG = 2
(IN25) GRASS(1,4);
                                                                               6
                                                       2
                                                           2
THE CHERNPOLYNOMIAL OF GRASS( 1 , 4 ) IS: (4 C2Q + 12 C1Q C2Q - 13 C1Q C2Q + 4 C1Q ) T
                                       5
                                                 2
                                                         2
                                                                                 3 3
                                                                                                     2
                        3
                                     5
  + (28 C1Q C2Q - 28 C1Q C2C + 18 C1Q ) T + (4 C2Q - 7 C1Q C2Q + 14 C1Q ) T + 15 C1Q T + (C2Q + 11 C1Q ) T + 5 C1
THE RELATIONS OF THE CHERNCLASSES OF Q ARE:
   2
          2
C2Q - 3 C1Q C2Q + C1Q = 0
         2
                 3
                            5
3 C10 C20 - 4 C10 C20 + C10
                              = 8
             2 2
     3
                           4
                                     6
- C2Q + 6 C1Q C2Q - 5 C1Q C2Q + C1Q = 8
                                                              5
                                                                       2
                                                                                 6
                              2
                                        3
GRMMR = - 56 - 13 D 55 - 78 D 54 - 286 D 53 - 715 D 52 - 1287 D 51 + DEG - 1716 D
    12
   8
                          2
                                     3
                                                          5
      = S6 + 12 D S5 + 66 D S4 + 228 D S3 + 495 D S2 + 792 D S1 + 924 D
RAM
   11
  8
                   2
                             3
                                                   5
 RAM
      = D S5 + 11 D S4 + 55 D S3 + 165 D S2 + 330 D S1 + 462 D
   10
 = 8
        2
                   3
                                        5
                                                   6
 RAM = D S4 + 10 D S3 + 45 D S2 + 120 D S1 + 210 D
   9
 = 8
        3
                            5
 RAM = D S3 + 9 D S2 + 36 D S1 + 84 D
   8
 = 2
 DEG = 5
```

```
(IN26) GRASS(1,5);
                                                      2
                                                          3
THE CHERNPOLYNOMIAL OF GRASS( 1 , 5 ) IS: (9 C20 - 18 C10 C20 + 42 C10 C20 - 26 C10 C20 + 5 C10 ) T
                                                   3
                                                            2
                                                                2
 + (75 C10 C20 - 68 C10 C20 + 15 C10 ) T + (- 6 C20 + 78 C10 C20 - 66 C10 C20 + 25 C10 ) T
                                                2
                                                        2
                                    5
                                       5
             2
+ (30 C10 C20 - 30 C10 C20 + 30 C10 ) T + (7 C20 - 2 C10 C20 + 31 C10 ) T + (6 C10 C20 + 26 C10 ) T + (2 C20 + 16
 + 6 C1Q T + 1
THE RELATIONS OF THE CHERNCLASSES OF Q ARE:
                 3
 C10 C20 - 4 C10 C20 + C10
            2
                                     6
     3
                2
- C2Q + 6 C1Q C2Q - 5 C1Q C2Q + C1Q = 8
                    3
                                 5
          3
                         2
    C10 C20 + 10 C10 C20 - 6 C10 C20 + C10 = 0
                          4
                               2
                                       6
            2 3
C2Q - 18 C1Q C2Q + 15 C1Q C2Q - 7 C1Q C2Q + C1Q
                              2
       = - S8 - 17 D S7 - 136 D S6 - 688 D S5 - 2388 D S4 - 6188 D S3 - 12376 D S2 - 19448 D S1 + DEG - 24318 D
    16
   A
                           2
      = S8 + 16 D S7 + 120 D S6 + 560 D S5 + 1820 D S4 + 4368 D S3 + 8008 D S2 + 11440 D S1 + 12870 D
RAM
  15
RRM = D S7 + 15 D S6 + 105 D S5 + 455 D S4 + 1365 D S3 + 3003 D S2 + 5005 D S1 + 6435 D
   14
. 8
      = D S6 + 14 D S5 + 91 D S4 + 364 D S3 + 1001 D S2 + 2002 D S1 + 3003 D
ROM
   13
 A
      = D S5 + 13 D S4 + 78 D S3 + 286 D S2 + 715 D S1 + 1287 D
RAM
   12
= 6
DEG = 14
 (IN27) GRASS(2,5);
                                                                      3
                                                                          2
                                                              2
THE CHERNPOLYNOMIAL OF GRASS( 2 , 5 ) IS: (8 C30 - 24 C10 C20 C30 + 4 C10 C30 + 52 C10 C20 C30 - 36 C10 C20 C30
                                  3
                                      3
                                               5
                                                    2
 + 7 C1Q C3Q - 8 C1Q C2Q - 28 C1Q C2Q + 38 C1Q C2Q - 21 C1Q C2Q + 4 C1Q ) T
  + (- 22 C10 C30 + 56 C10 C20 C30 - 12 C10 C20 C30 + 2 C10 C30 - 4 C20 - 48 C10 C20 + 85 C10 C20 - 57 C10 C20
                                      2
  + 14 C1Q ) T + (- 38 C1Q C3Q + 24 C2Q C3Q + 36 C1Q C2Q C3Q + 2 C1Q C3Q - 24 C1Q C2Q + 78 C1Q C2Q - 78 C1Q C2Q +
                                                                               6
    + (- 15 C30 + 30 C10 C20 C30 + 26 C10 C30 + 43 C10 C20 - 84 C10 C20 + 43 C10 ) T
                                    3
  + (41 C10 C30 + 22 C10 C20 - 63 C10 C20 + 58 C10 ) T + (24 C10 C30 + 3 C20 - 27 C10 C20 + 45 C10 ) T
                                          2 2
  + (6 C3Q - 6 C1Q C2Q + 32 C1Q ) T + 17 C1Q T + 6 C1Q T + 1
 THE RELATIONS OF THE CHERNCLASSES OF Q ARE:
               2
                       2
 2 C10 C30 + C20 - 3 C10 C20 + C10 = 8
                                  2
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 - 2 C2Q C3Q + 3 C1Q C3Q + 3 C1Q C2Q - 4 C1Q C2Q + C1Q
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1:302 - 6 010 020 030 + 4 0103 030 - 0203 + 6 01020202 - 5 0104 020 + 010 = 8 2 2. 2 4 3 3 2 5 7 1 C 1 Q C 3 Q + 3 C 2 Q C 3 Q - 1 2 C 1 Q C 2 Q C 3 Q + 5 C 1 Q C 3 Q - 4 C 1 Q C 2 Q + 1 8 C 1 Q C 2 Q - 6 C 1 Q C 2 Q + C 1 Q = 8 2 2 2 2 2 3 5 4 2 3 4 2 6 - 3 C20 C30 + 6 C10 C30 + 12 C10 C20 C30 - 28 C10 C20 C30 + 6 C10 C30 + C20 - 18 C10 C20 + 15 C10 C20 - 7 C10 C + C10
 3
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 - 28
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 + 21 CIQ C2Q - 8 CIQ C2Q + CIQ 2 3 4 5 6 7 8 - S9 - 19 D S8 - 171 D S7 - 969 D S6 - 3876 D S5 - 11628 D S4 - 27132 D S3 - 58388 D S2 - 75582 D S1 + DEG AHHA 18 92 2 = S9 + 18 D S8 + 153 D S7 + 816 D S6 + 3060 D S5 + 8568 D S4 + 18564 D S3 + 31824 D S2 + 43758 D S1 + 48620 D g tan 17 4 EG = 42(IN28) GRASS(2,6);

ORE capacity exceeded (while requesting LIST space)

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